# PCFGs: Parsing & Evaluation

LING 571 — Deep Processing Techniques for NLP Shane Steinert-Threlkeld

#### Announcements

- HW2 due tonight at 11:59pm
  - readme.{txt|pdf}
    - Separate upload to Canvas
    - **NOT in** hw2.tar.gz
  - Run check\_hw2.sh before submitting! (Also: tar -tf hw2.tar.gz to preview.)
    - Flat structure; just files, no directories, inside tar-ball
    - Include *only* the files we ask for, not more
- Start symbol: either "%start S" or first nonterminal
  - NB: needs to be readable by nltk's grammar loading methods
- Use nltk.data.load: best to use "file:path/to/grammar.cfg" as argument
  - Docs: <a href="https://www.nltk.org/api/nltk.data.html#nltk.data.load">https://www.nltk.org/api/nltk.data.html#nltk.data.load</a>
- See hw2 slides as well on website for above points

```
(base) [shanest@patas ref]$ tar -tf hw1.tar.gz
hw1.py
hw1_parse.out
run_hw1.sh
```

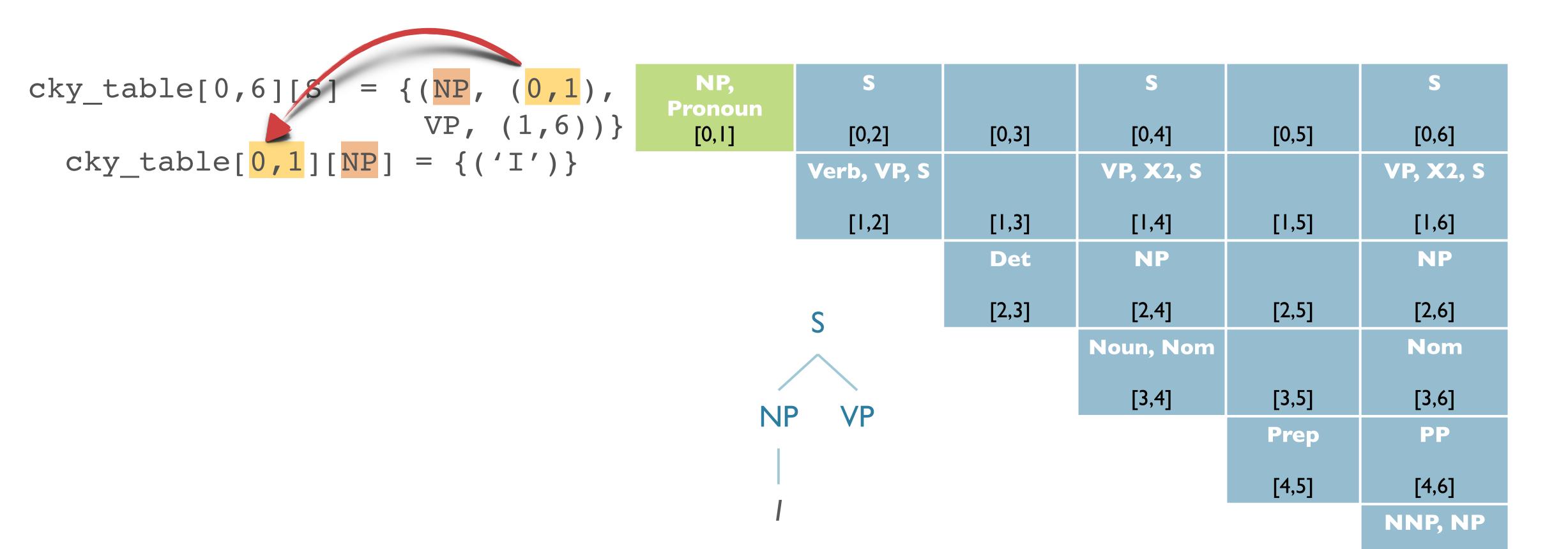
## Roadmap

- CKY + back-pointers example
- PCFGs
- PCFG Parsing (PCKY)
- Inducing a PCFG
- Evaluation
- [Earley parsing]
- HW3 + collaboration

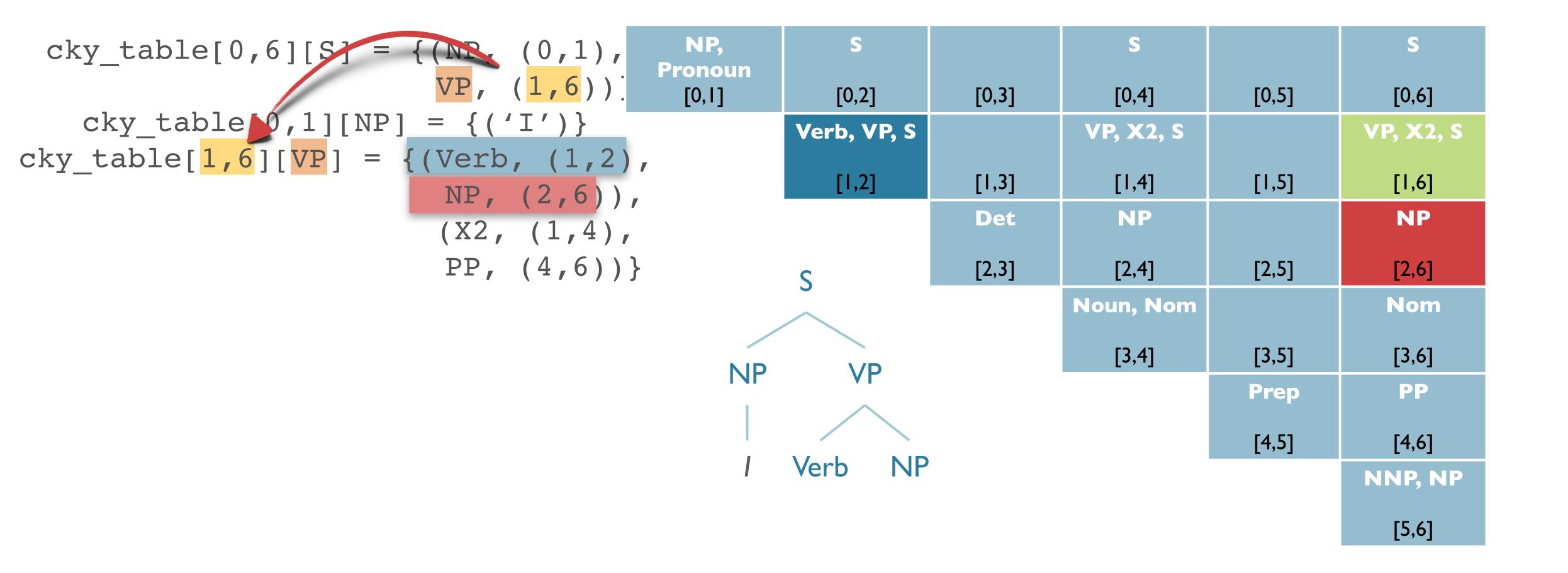
# CKY + Back-pointers Example

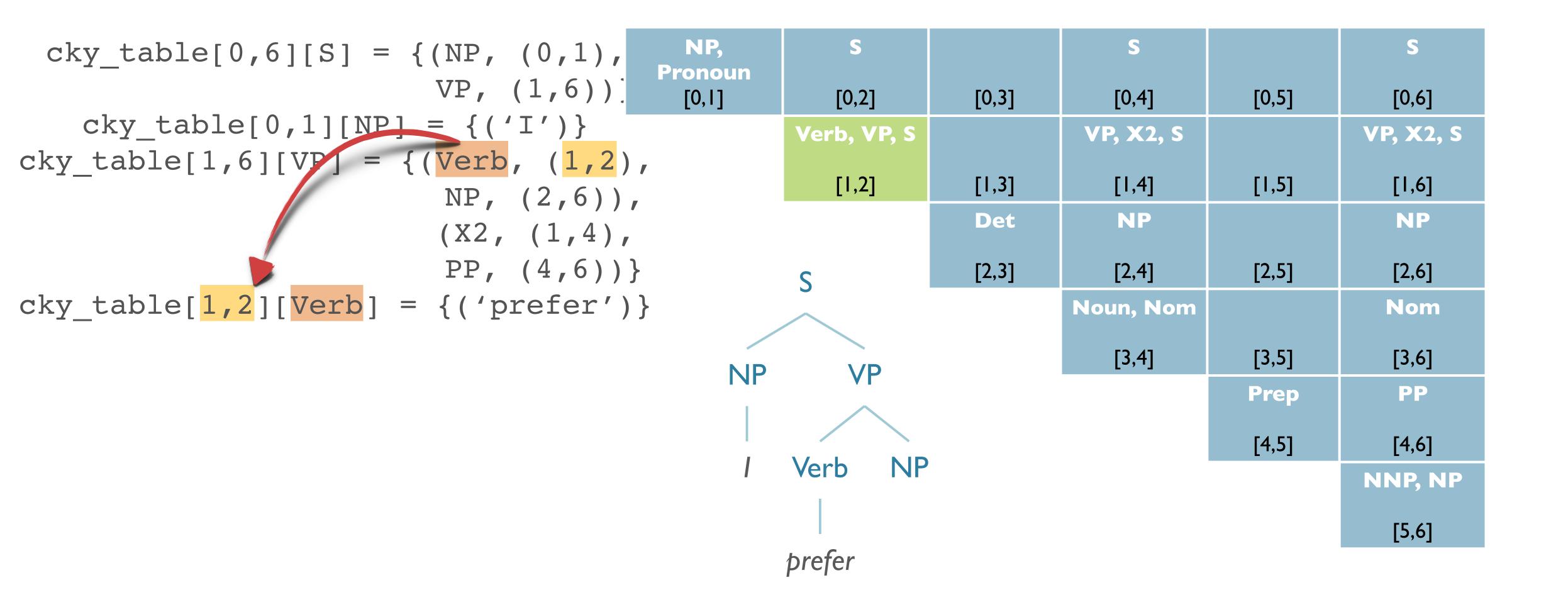
 $cky\_table[0,6][S] = \{ (NP, (0,1), VP, (1,6)) \}$ NP, **Pronoun** [0,1] [0,2] [0,3] [0,4] [0,5] [0,6] Verb, VP, S **VP, X2, S VP, X2, S** [1,2] [1,3] [1,4] [1,5] [1,6] NP Det NP [2,3] [2,4] [2,5] [2,6] Noun, Nom Nom [3,4] [3,5] [3,6] NP VP PP Prep [4,6] [4,5] NNP, NP [5,6]

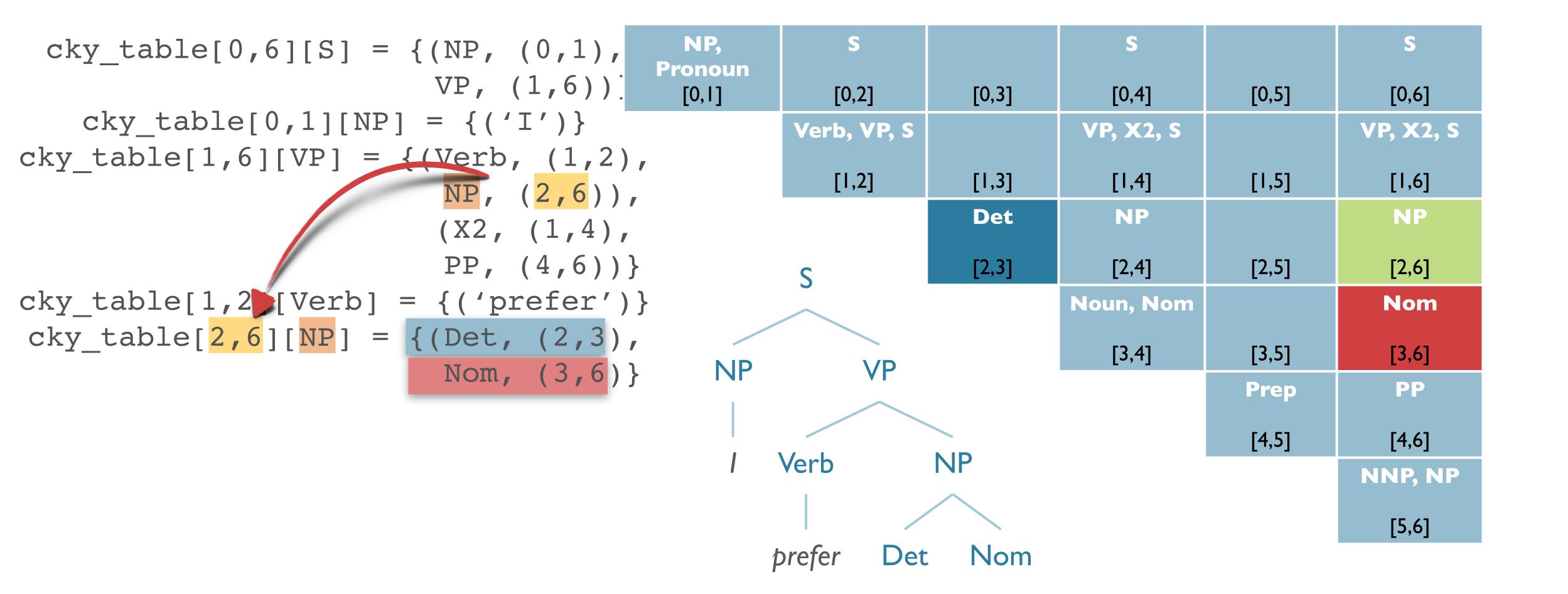
prefer a flight on TWA

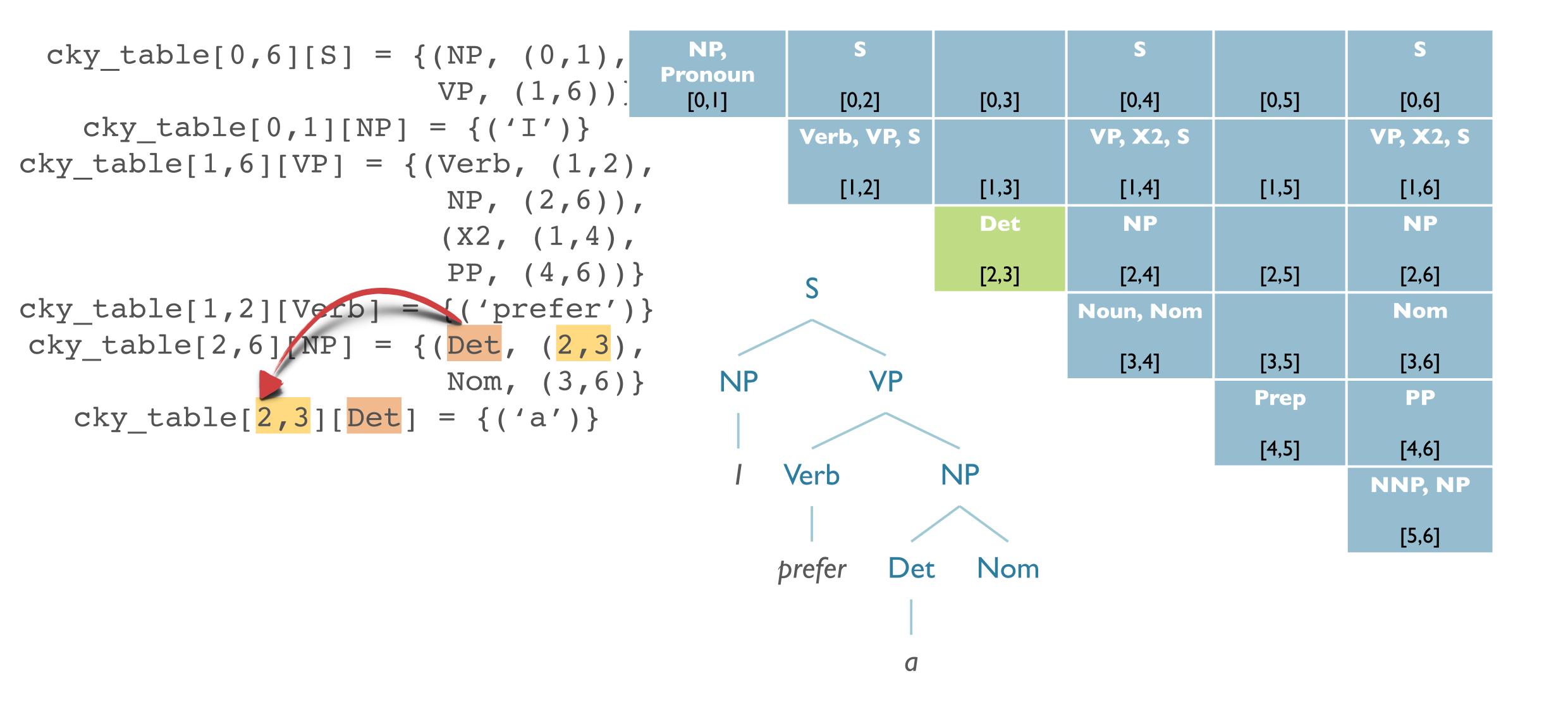


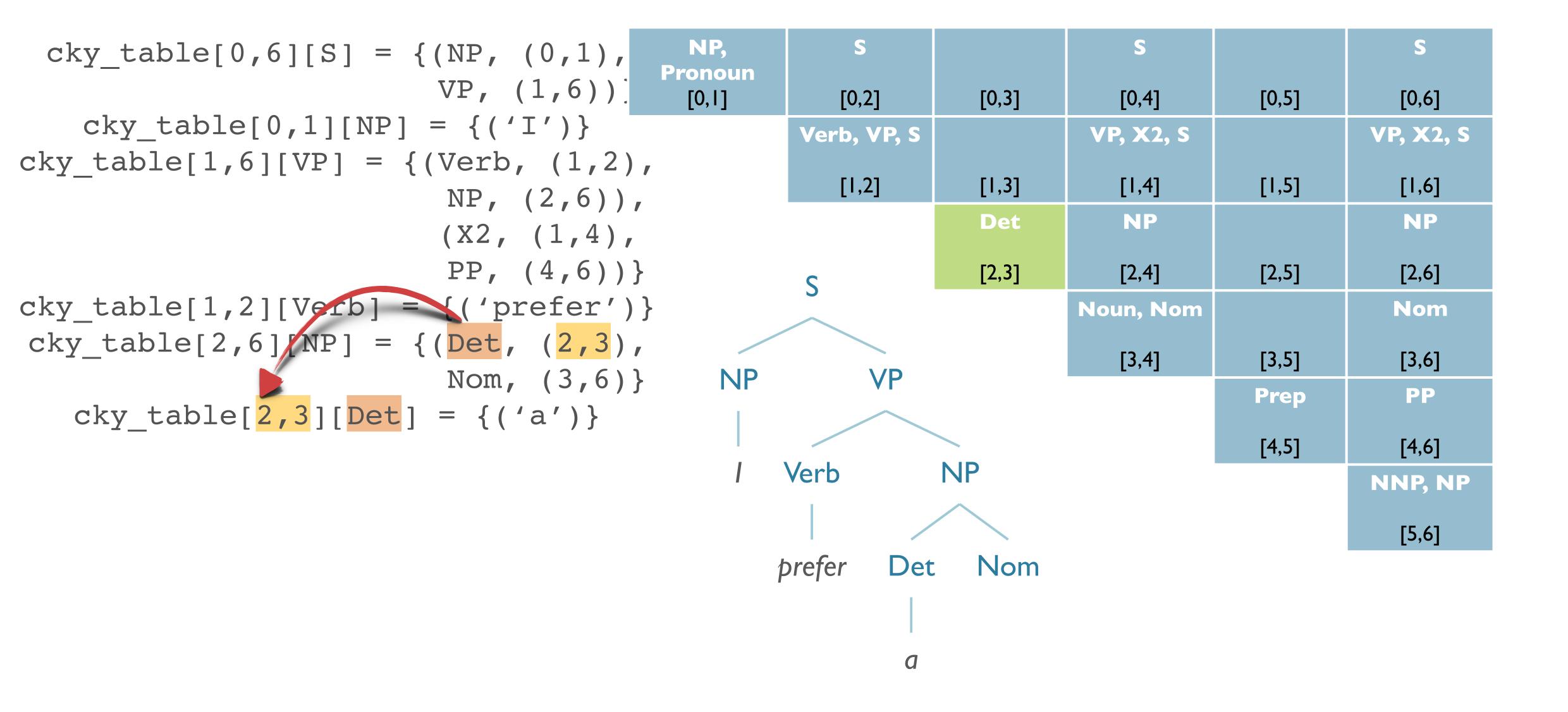
[5,6]



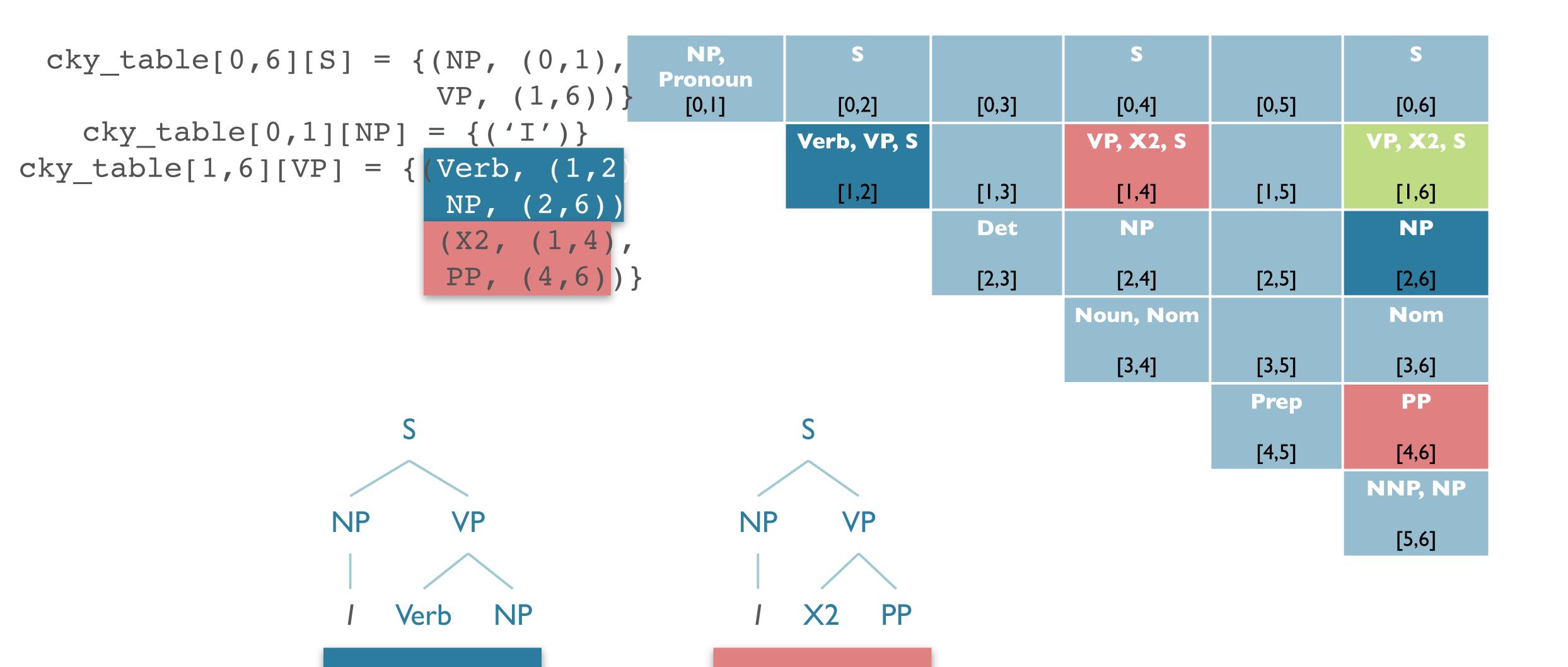








TWA



## Probabilistic Context-Free Grammars

## Probabilistic Context-free Grammars: Roadmap

Motivation: Ambiguity

Approach:

Definition

Disambiguation

Parsing

Evaluation

Enhancements

### Motivation

What about ambiguity?

Current algorithm can *represent* it

...can't resolve it.

# Probabilistic Parsing

- Provides strategy for solving disambiguation problem
  - Compute the probability of all analyses
  - Select the most probable

- Employed in language modeling for speech recognition
  - N-gram grammars predict words, constrain search
  - Also, constrain generation, translation

a set of non-terminal symbols (or variables)

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N	a set of <b>non-terminal symbols</b> (or <b>variables</b> )		
Σ	a set of <b>terminal symbols</b> (disjoint from <b>N</b> )		
R	a set of rules of productions, each of the form $A \to \beta[p]$ , where $A$ is a non-terminal where $A$ is a non-terminal, $\beta$ is a string of symbols from the infinite set of strings $(\Sigma \cup N) *$ and $p$ is a number between $0$ and $1$ expressing $P(\beta   A)$		

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S	a designated <b>start symbol</b>	

#### PCFGs

- Augment each production with probability that LHS will be expanded as RHS
  - $\bullet P(A \rightarrow \beta)$
  - $P(A \rightarrow \beta | A)$
  - $\bullet$   $P(\beta|A)$
  - P(RHS | LHS)
- NB: the first is often used; but the latter are what's really meant.

#### PCFGs

Sum over all possible expansions is 1

$$\sum_{\beta} P(A \to \beta) = 1$$

- A PCFG is consistent if sum of probabilities of all sentences in language is **1** 
  - Recursive rules often yield inconsistent grammars (<u>Booth & Thompson, 1973</u>)

# Example PCFG: Augmented $\mathcal{L}_1$

<u> </u>		
Grammar		Lexicon
$S \rightarrow NP VP$	[.80]	Det → that [.10]   a [.30]   the [.60]
$S \rightarrow Aux NP VP$	[.15]	<i>Noun</i> → <i>book</i> [.10]   <i>flight</i> [.30]   <i>meal</i> [.15]   <i>money</i> [0.5]
$S \rightarrow VP$	[.05]	l flights [0.40] l dinner [.10]
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NP → Proper-Noun	[.30]	<i>Pronoun</i> → $I[.40]   she[.05]   me[.15]   you[.40]$
NP → Det Nominal	[.20]	Proper-Noun → Houston [.60]   NWA [.40]
NP → Nominal	[.15]	<i>Aux</i> → <i>does</i> [.60] I <i>can</i> [.40]
Nominal → Noun	[.75]	Preposition $\rightarrow$ from [.30]   to [.30]   on [.20]   near [.15]
Nominal → Nominal Noun	[.20]	l <i>through</i> [.05]
Nominal → Nominal PP	[.05]	
VP → Verb	[.35]	
VP → Verb NP	[.20]	
VP → Verb NP PP	[.10]	
VP → Verb PP	[.15]	
VP → Verb NP NP	[.05]	
$VP \rightarrow VP PP$	[.15]	
PP → Preposition NP	[1.0]	

# Example PCFG: Augmented £1

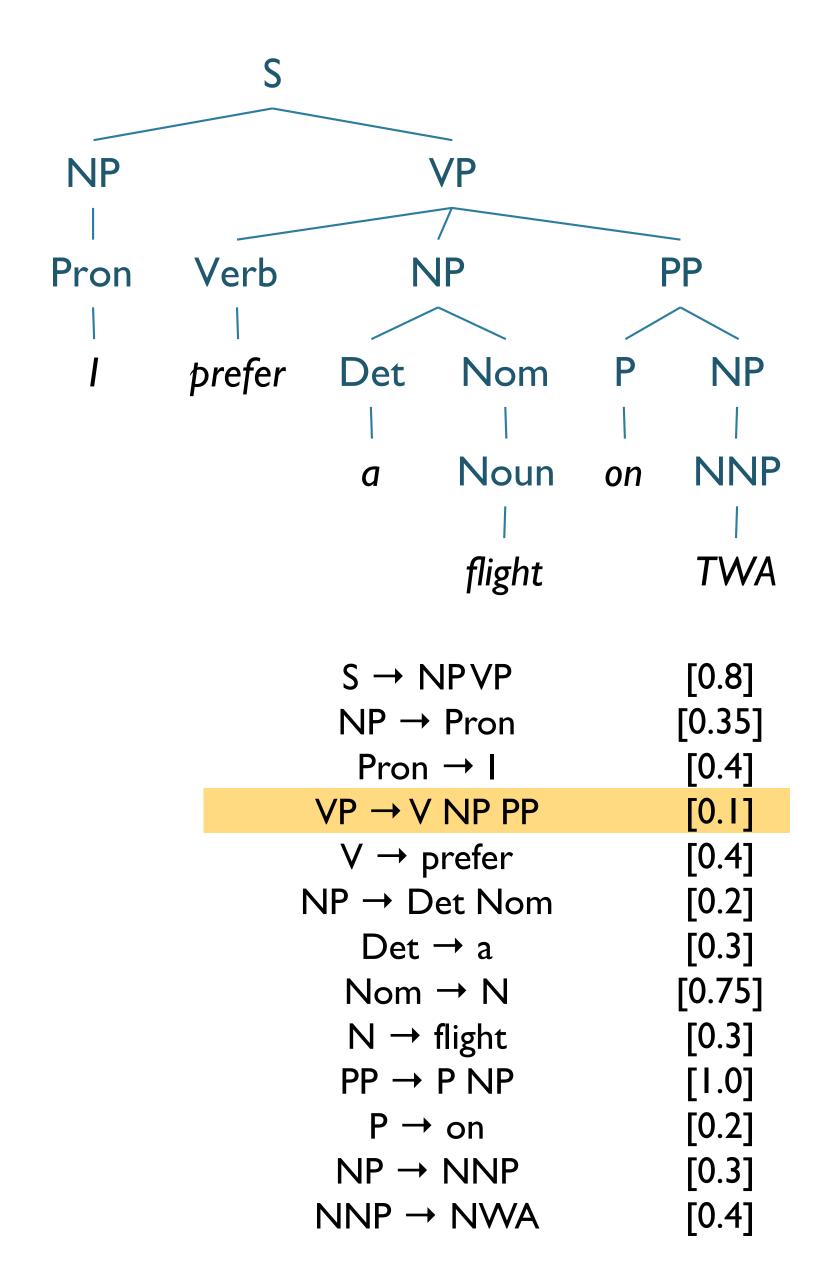
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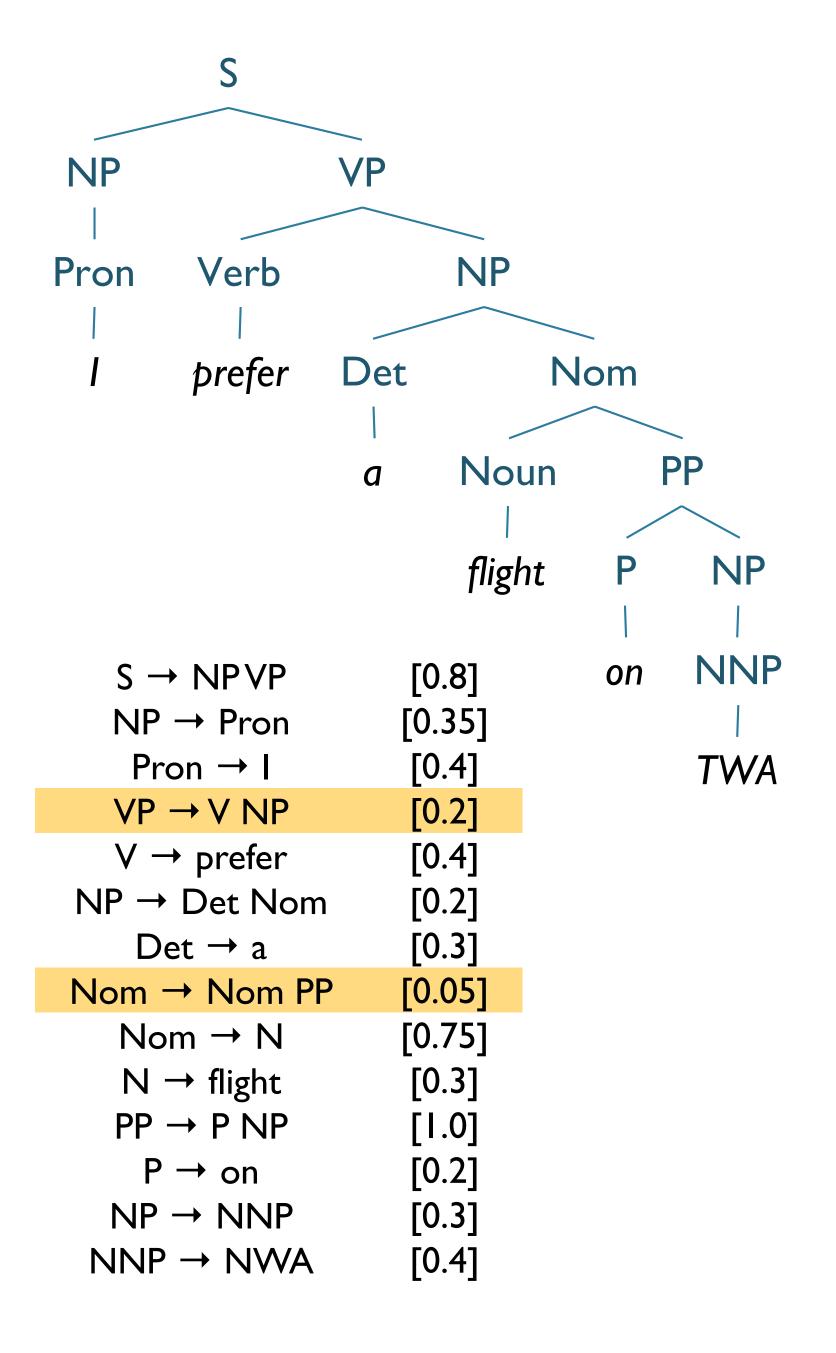
# Disambiguation

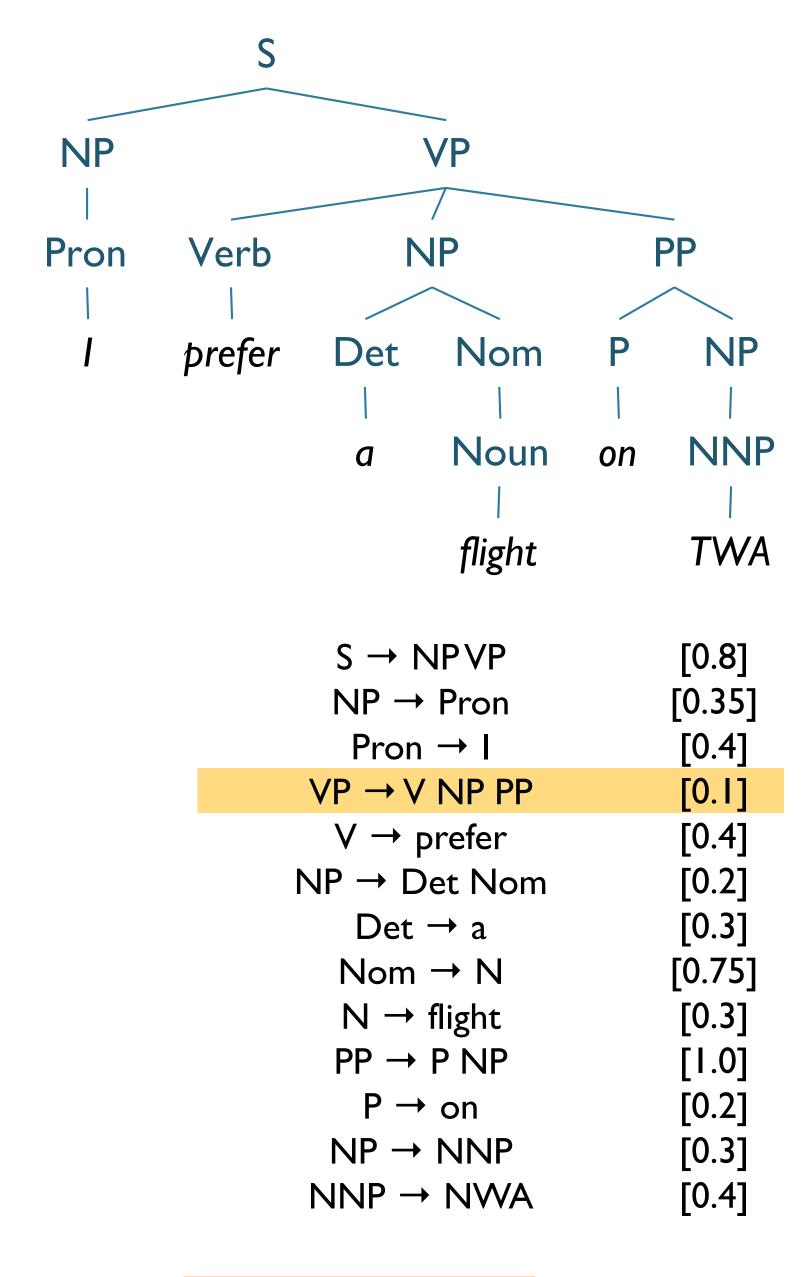
- A PCFG assigns probability to each parse tree T for input S
- Probability of T: product of all rules used to derive T

$$P(T,S) = \prod_{i=1}^{n} P(RHS_i | LHS_i)$$

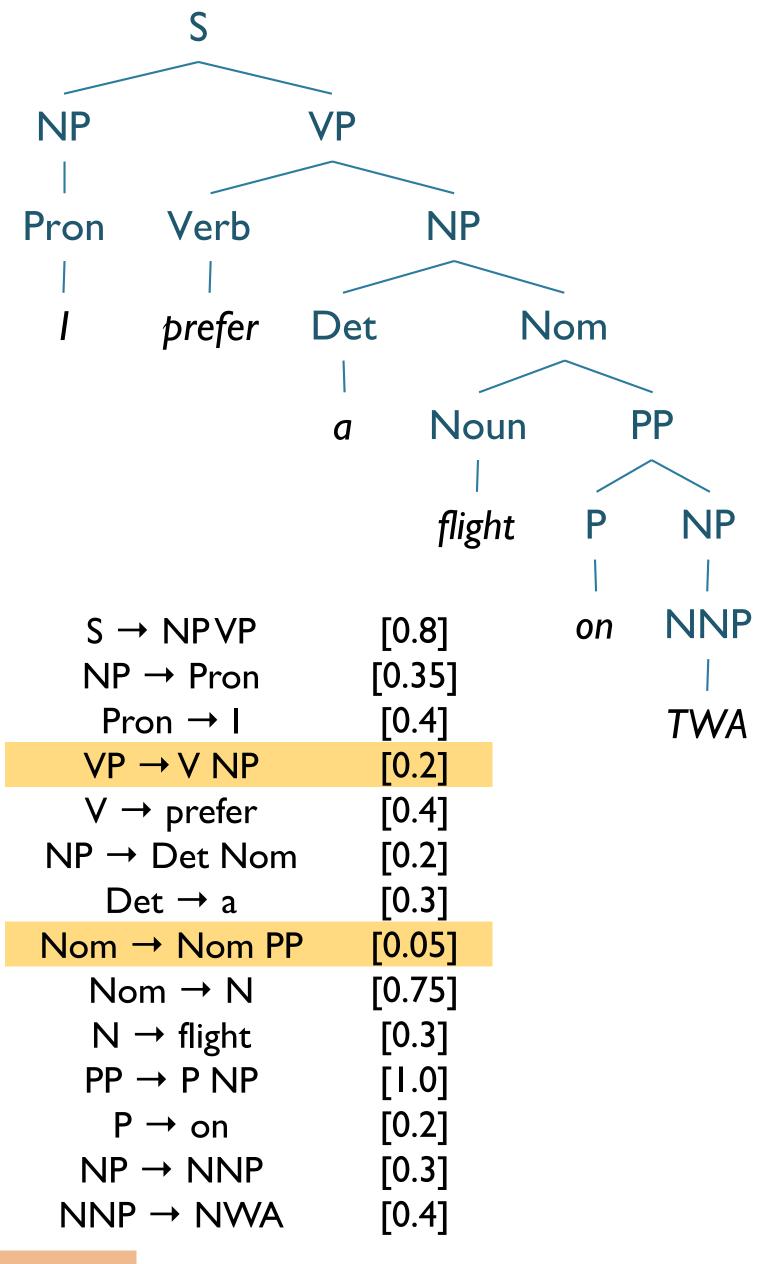
$$P(T, S) = P(T)P(S | T) = P(T)$$







 $\sim 1.452 \times 10^{-6}$ 



 $\sim 1.452 \times 10^{-7}$ 

# Parsing Problem for PCFGs

• Select **T** such that **(s.t.)** 

$$\hat{T}(S) = \underset{T \text{ s.t. } S = yield(T)}{\operatorname{argmax}} P(T)$$

- String of words S is yield of parse tree
- Select the tree Tthat maximizes the probability of the parse

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- PCFGs are able to give probability of entire string without as bad sparsity
- Model probability of syntactically valid sentences
  - Not just probability of sequence of words

# PCFGs: Parsing

### Probabilistic CKY (PCKY)

- Like regular CKY
  - Assumes grammar in Chomsky Normal Form (CNF)
    - $\bullet$   $A \rightarrow B C$
    - $\bullet$   $A \rightarrow W$
  - Represent input with indices b/t words:
    - Book 1 that 2 flight 3 through 4 Houston 5

### Probabilistic CKY (PCKY)

- For input string length n and non-terminals V
  - Cell [ i, j, A ] in ( n+1 ) × ( n+1 ) × V matrix
  - Contains probability that A spans [i, j]

```
function Probabilistic-CKY-PARSE(words, grammar) returns most probable parse and its probability
for j ← from 1 to LENGTH(words) do
for all \{A \mid A \rightarrow words[j] \in grammar\}
      table[i-1, j, A] \leftarrow P(A \rightarrow words[j])
 for i \leftarrow \text{from } j-2 \text{ downto } 0 \text{ do}
  for k \leftarrow i + 1 to j-1 do
  for all \{A \mid A \rightarrow B \ C \in grammar, \}
        and table[i, k, B] > 0 and table[k, j, C] > 0
  if (table[i, j, A] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]) then
      table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]
      back[i, j, A] \leftarrow \{k, B, C\}
  return Build_Tree(back[ 1, Length(words), S]), table[ 1,Length(words), S]
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### PCKY Grammar Segment

```
S \rightarrow NP VP  [0.80]
                                             Det \rightarrow the
                                                                [0.40]
NP \rightarrow Det N  [0.30]
                                              Det \rightarrow a
                                                            [0.40]
VP \rightarrow V NP \quad [0.20]
                                          V \rightarrow \text{includes} [0.05]
                                             N \rightarrow \text{meal}
                                                                [0.01]
                                             N → flight
                                                                [0.02]
```

 $S \rightarrow NP VP$  [0.80]  $NP \rightarrow Det N$  [0.30]  $VP \rightarrow V NP \quad [0.20]$ 

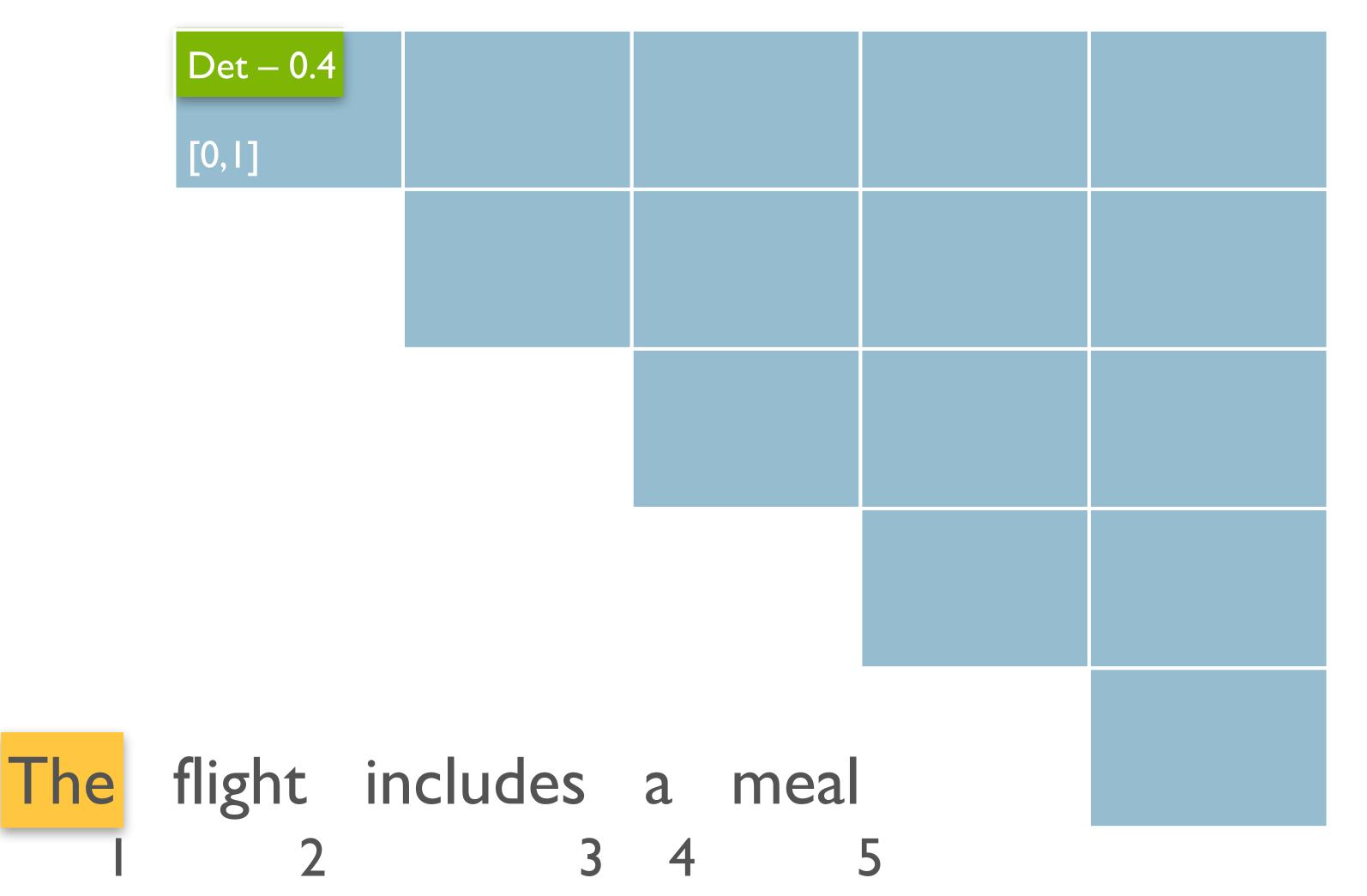


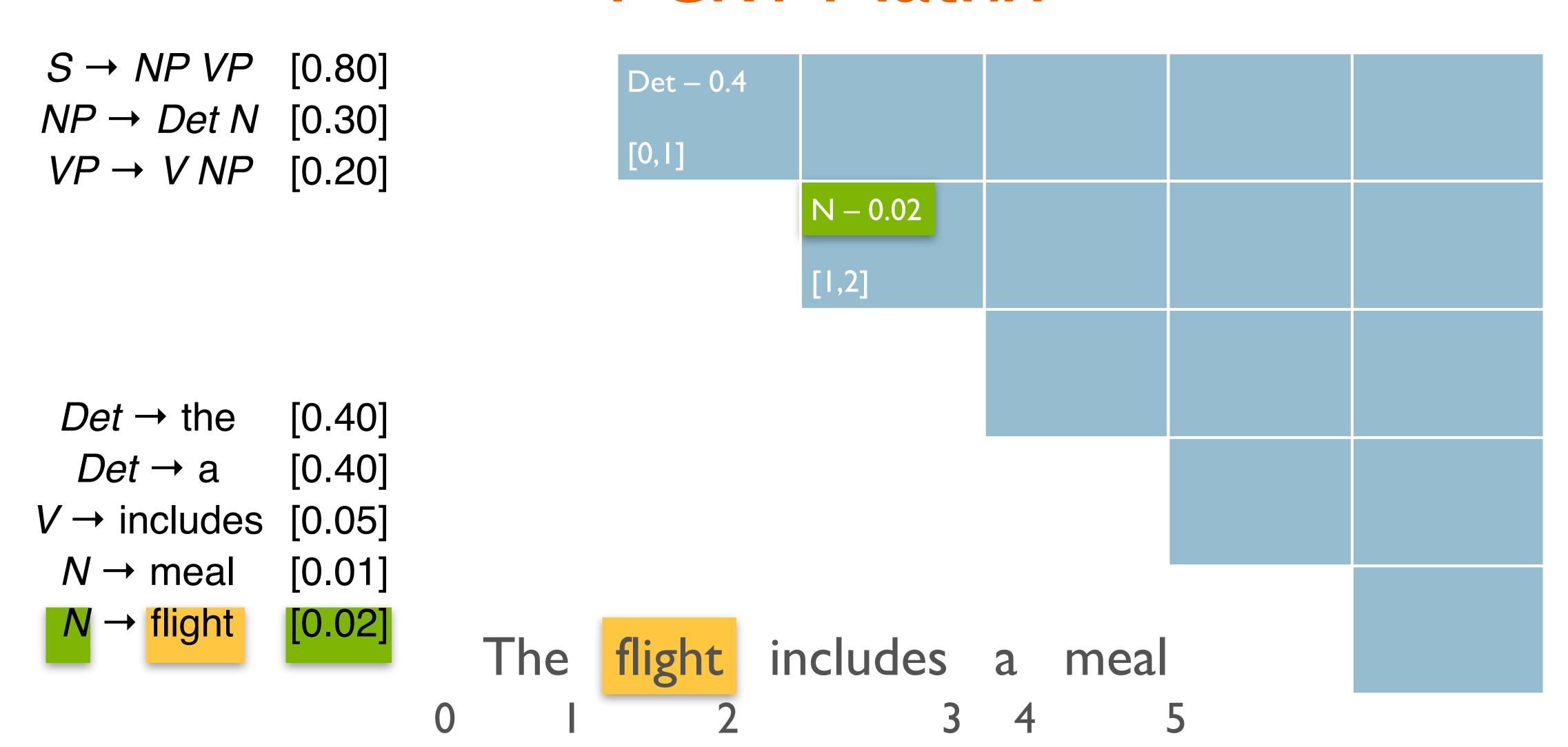
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 $V \rightarrow \text{includes} [0.05]$ 

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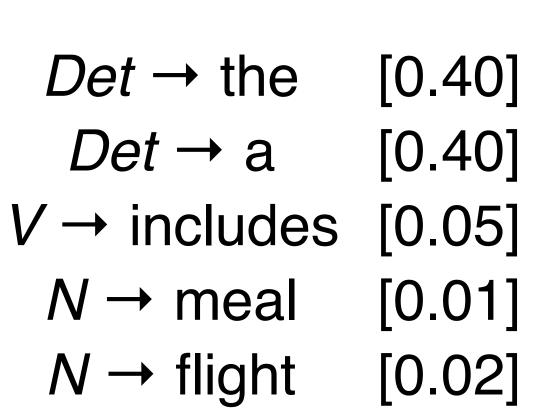
N → flight [0.02]

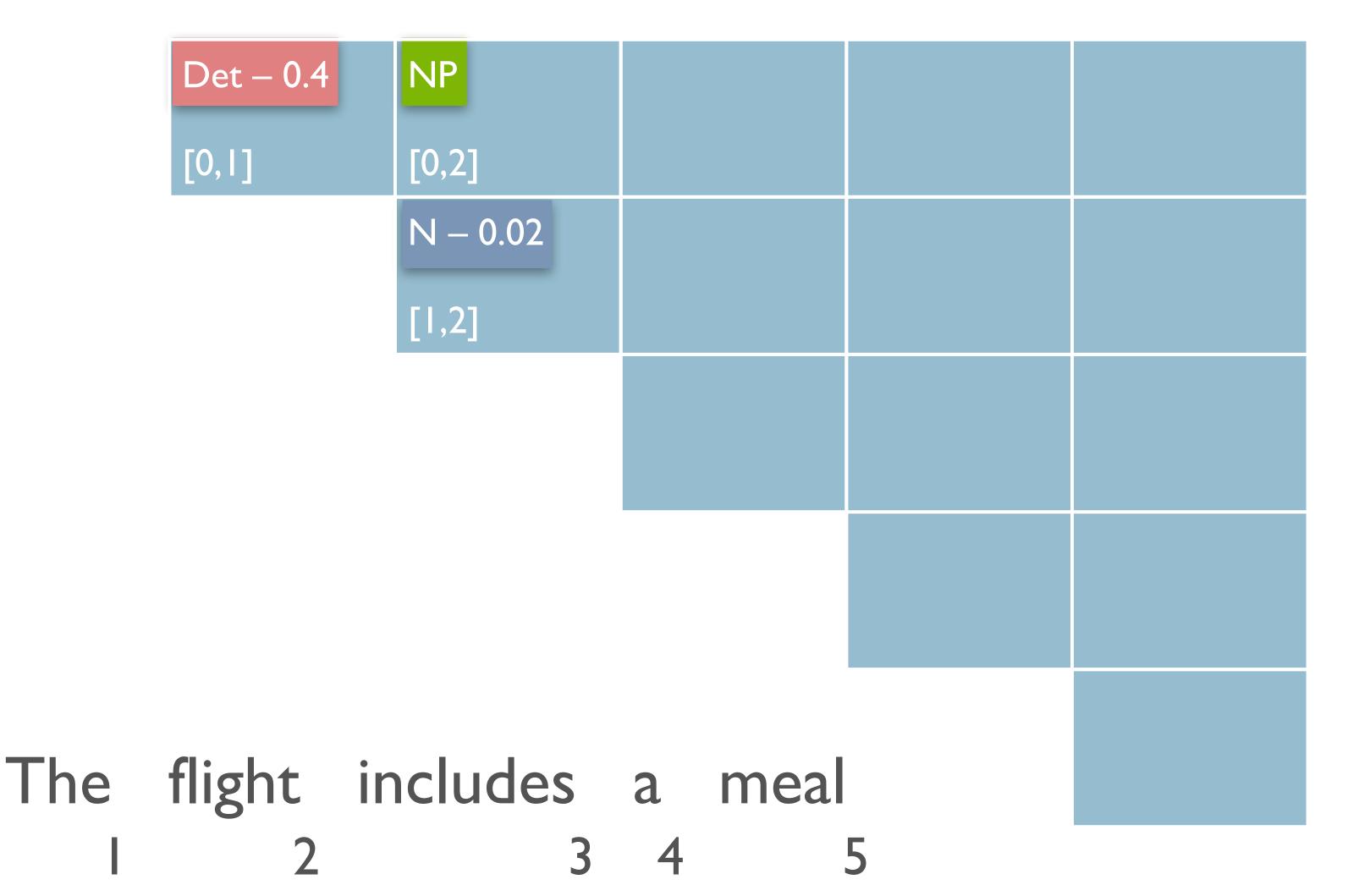




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```
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                    [0.40]
   Det \rightarrow a
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                    [0.01]
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                    [0.02]
```

```
Det – 0.4
                 NP
       [0,1]
                 [0,2]
                 N - 0.02
                 [1,2]
The flight includes a meal
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 [0.80]  
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[0.02]

N → flight

```
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        [0,1]
                    [0,2]
                   N - 0.02
                   [1,2]
       P(N \rightarrow flight)
                 = 0.00024
The flight includes a meal
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                       [0.01]
  N \rightarrow \text{meal}
```

[0.02]

N → flight

```
NP - 0.0024
           Det – 0.4
           [0,1]
                        [0,2]
                        N - 0.02
                        [1,2]
         P(Det \rightarrow a)
         P(N \rightarrow flight)
P = 0.3 \cdot 0.4 \cdot 0.02 = 0.00024
  The flight includes a meal
```

$S \rightarrow NP VP$	[0.80]
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$Det \rightarrow a$	[0.40]
V → includes	[0.05]
N → meal	[0.01]
N → flight	[0.02]

The

Det - 0.4	NP - 0.0024			S - 2.304×10-8
[0,1]	[0,2]	[0,3]	[0,4]	[0,5]
	N - 0.02			
	[1,2]	[1,3]	[1,4]	[1,5]
		V — 0.05		VP - 1.2×10-5
		[2,3]	[2,4]	[2,5]
			Det – 0.4	NP - 0.0012
			[3,4]	[3,5]
			N - 0.0 I	
flight includes a meal			[4,5]	

# Inducing a PCFG

- Simplest way:
  - Use treebank of parsed sentences

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 Number of times a nonterminal is expanded by a given rule:  $Count(\alpha \rightarrow \beta)$ 

$$P(\alpha \to \beta \mid \alpha) = \frac{Count(\alpha \to \beta)}{\sum_{\gamma} Count(\alpha \to \gamma)} = \frac{Count(\alpha \to \beta)}{Count(\alpha)}$$

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- Number of times a nonterminal is expanded by a given rule:

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- Alternative: Learn probabilities by re-estimating
  - (Later)

### Probabilistic Parser Development Paradigm

	Train	Dev	Test
	Large	Small	Small/Med
Size	(eg.WSJ 2-21, 39,830 sentences)	(e.g.WSJ 22)	(e.g. WSJ, 23, 2,416 sentences)
Usage	Estimate rule probabilities	Tuning/Verification, Check for Overfit	Held Out, Final Evaluation

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- How can we tell how good a parse is?
  - Maximally strict: identical to 'gold standard'
  - Partial credit:
    - Constituents in output match those in reference
      - Same start point, end point, non-terminal symbol

#### Parseval

- How can we compute parse score from constituents?
- Multiple Measures:

```
# of correct constituents in hypothetical parse
 Labeled Recall (LR) =
                             # of total constituents in reference parse
                          # of correct constituents in hypothetical parse
Labeled Precision (LP) =
                            # of total consituents in hypothetical parse
```

#### Parseval

#### • F-measure:

- Combines precision and recall
- Let  $\beta \in \mathbb{R}$ ,  $\beta > 0$  that adjusts P vs. R s.t.

$$\beta \propto \frac{R}{P}$$

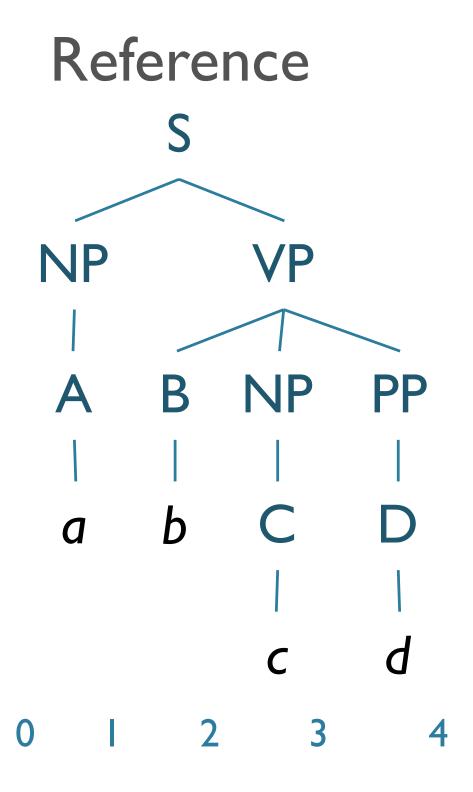
•  $F_{\beta}$ -measure is then:

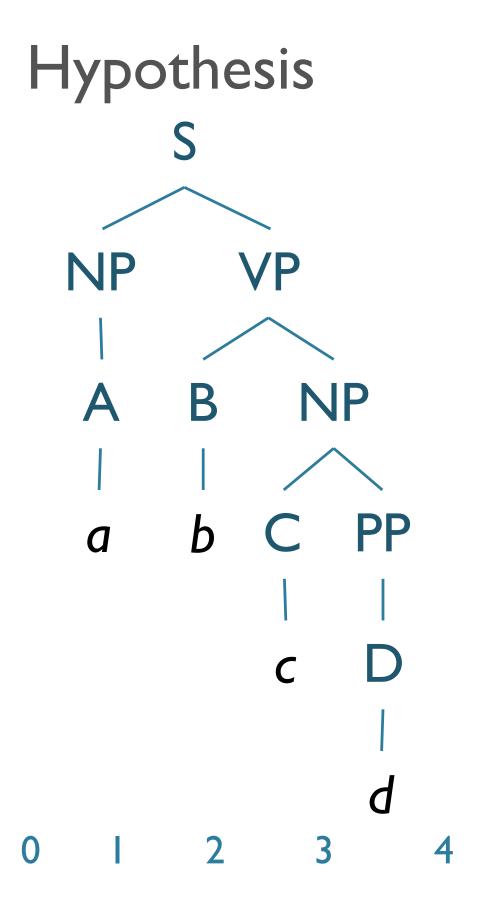
$$F_{\beta} = (1 + \beta^2) \cdot \frac{P \cdot R}{\beta^2 \cdot P + R}$$

With F1-measure as

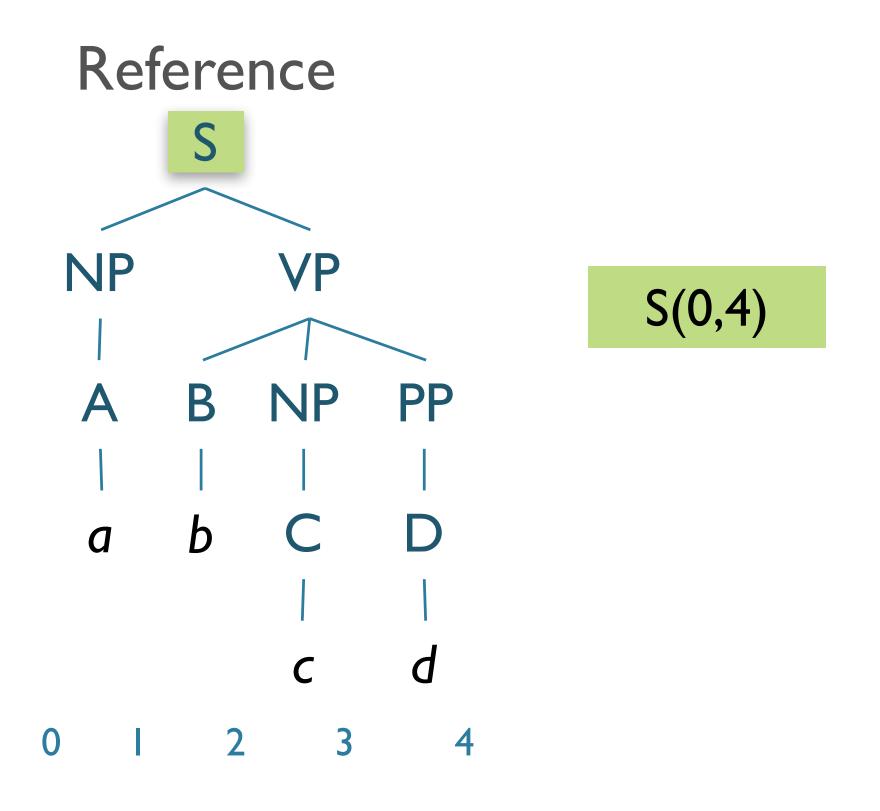
$$F_1 = \frac{2PR}{P + R}$$

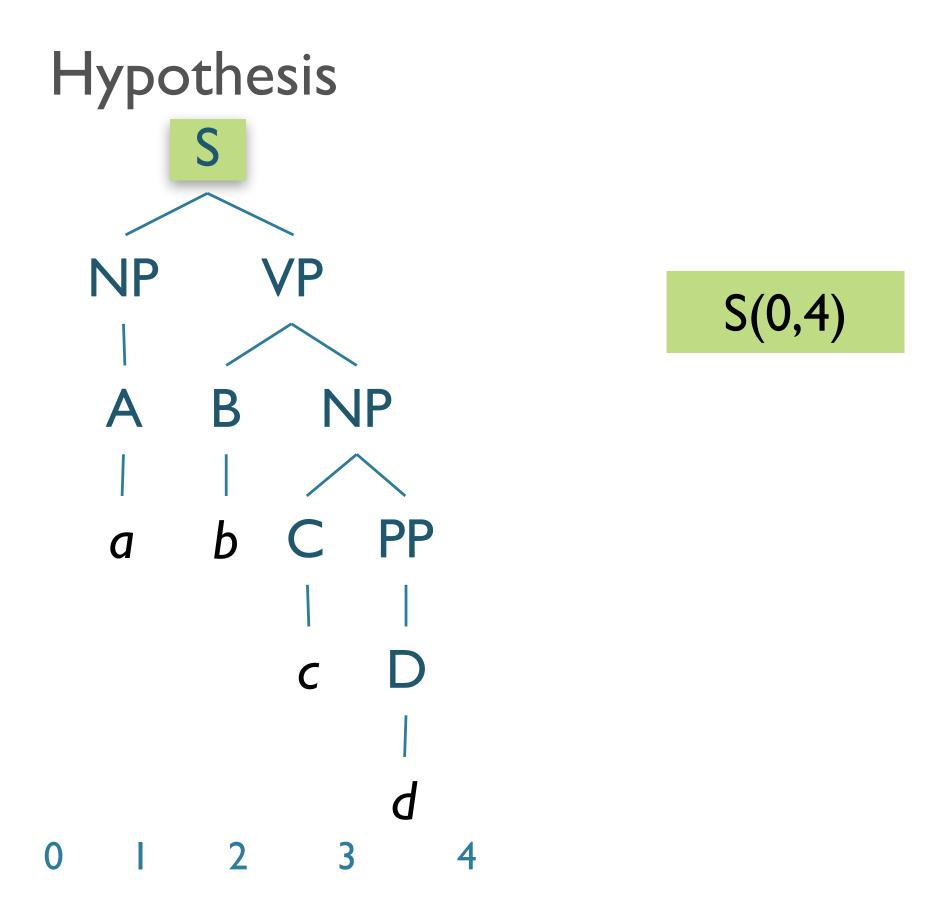
### Evaluation: Example



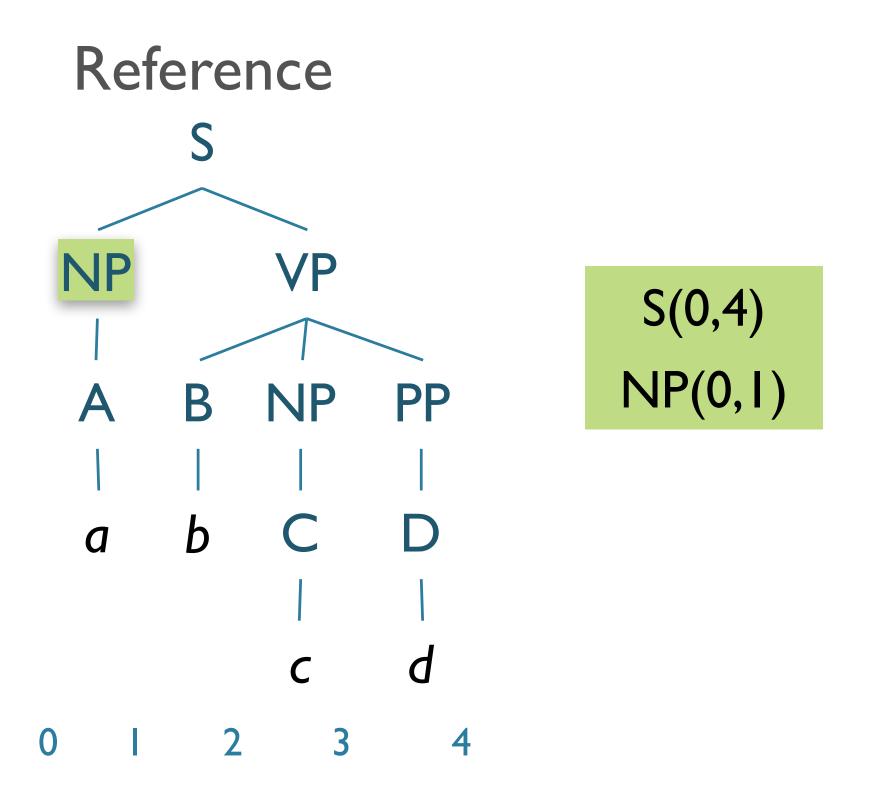


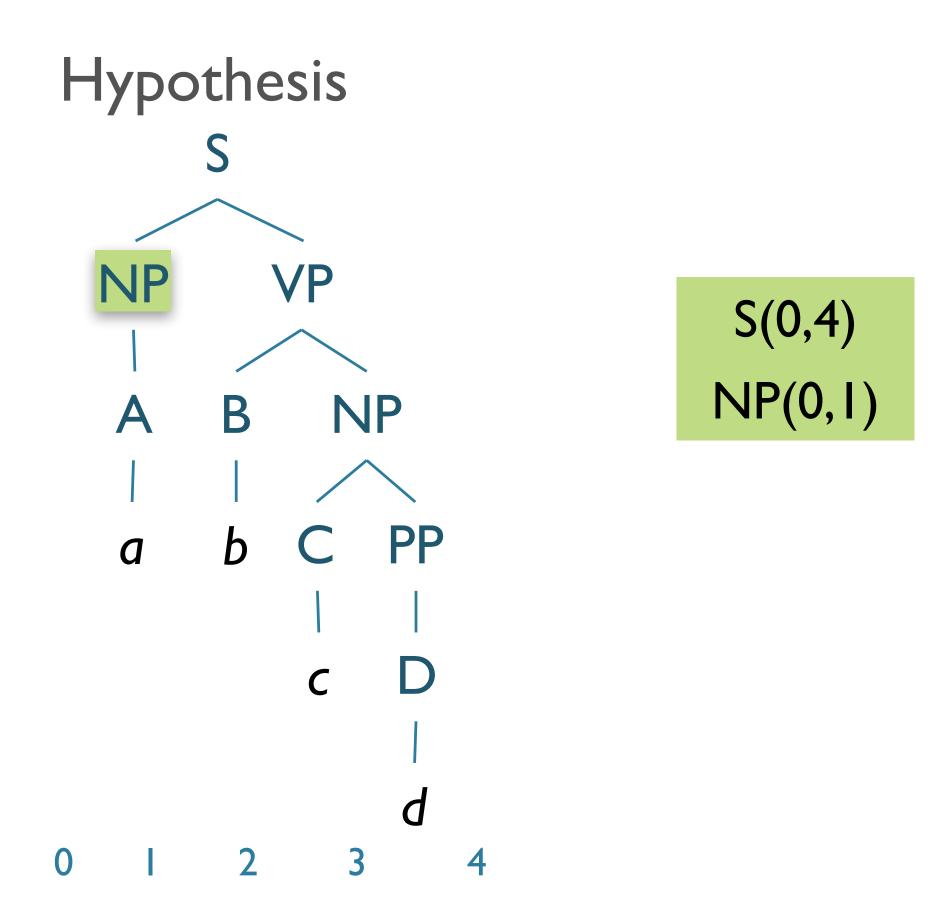
### Evaluation: Example

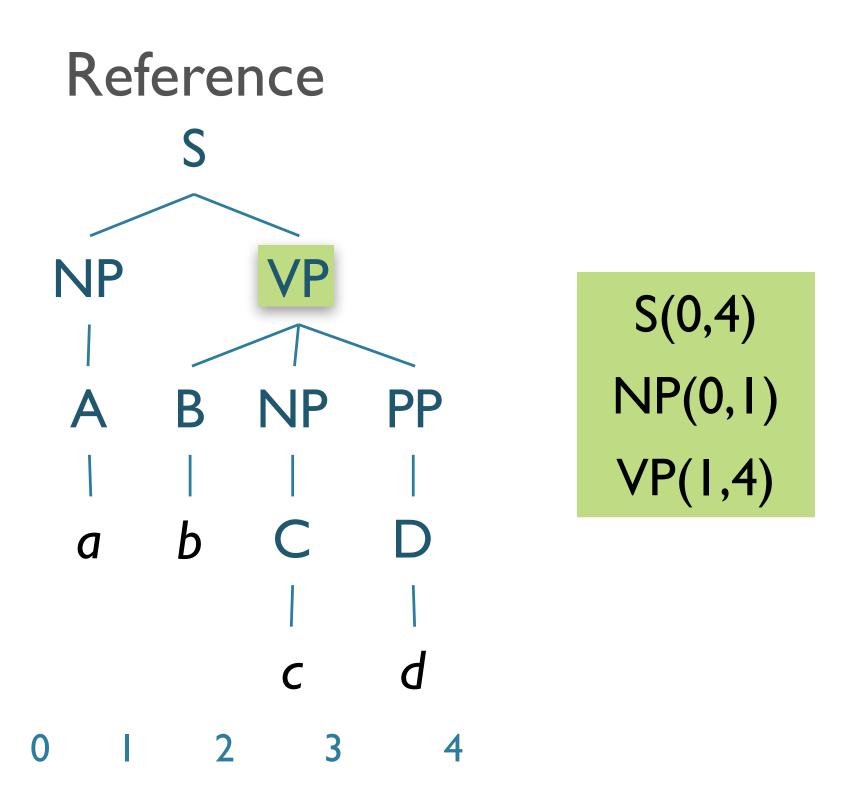


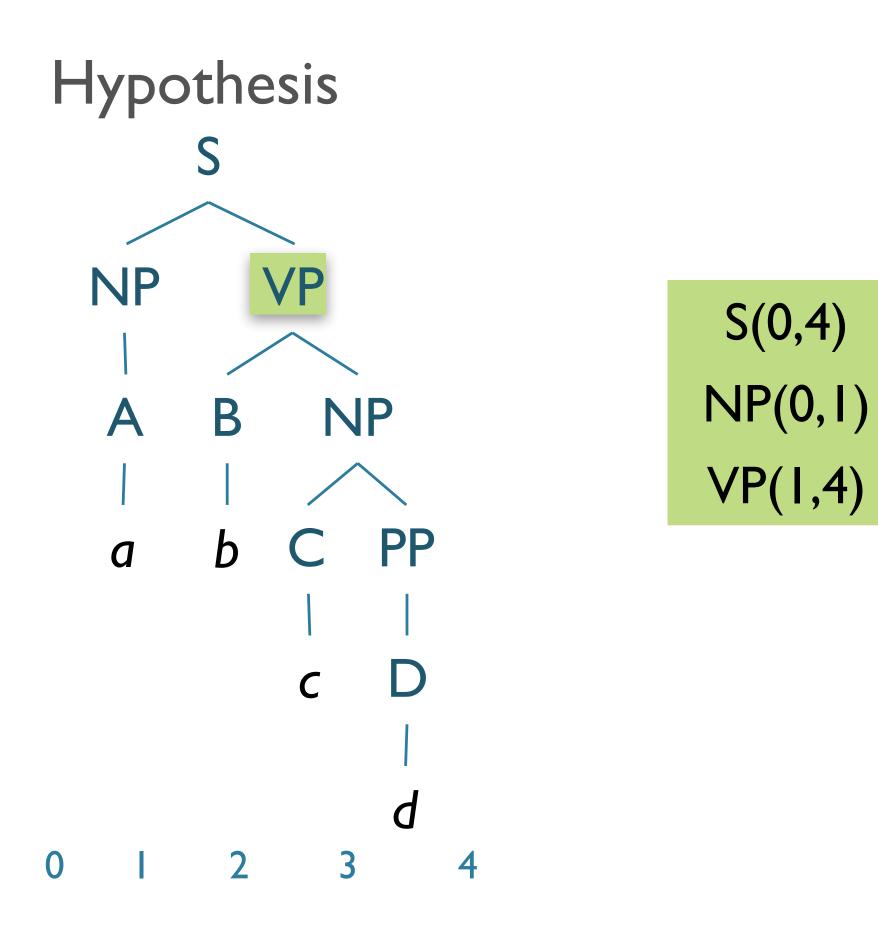


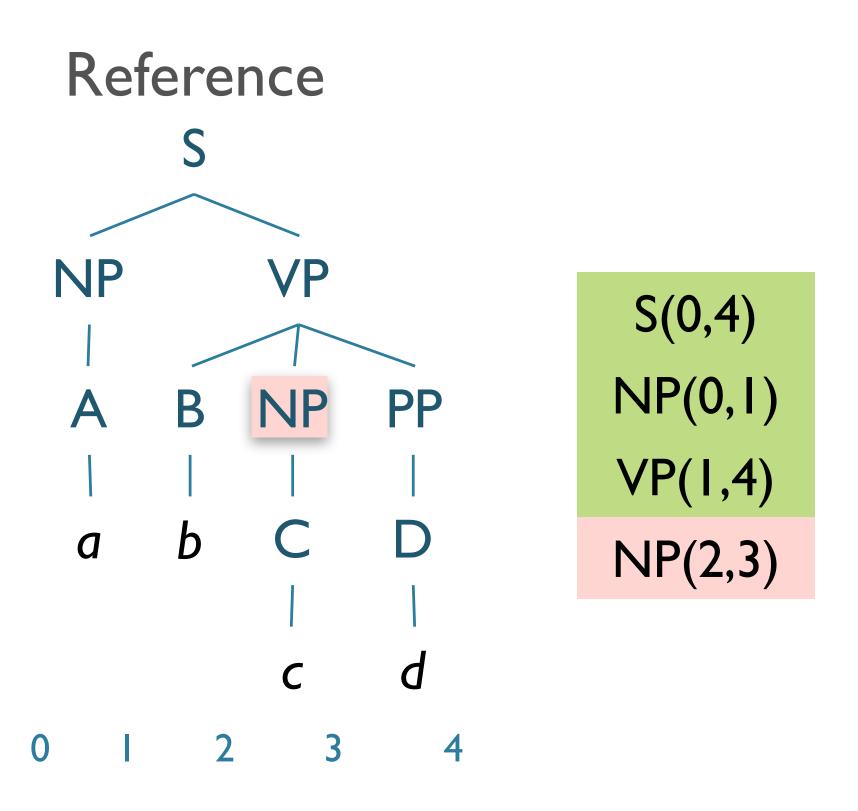
### Evaluation: Example

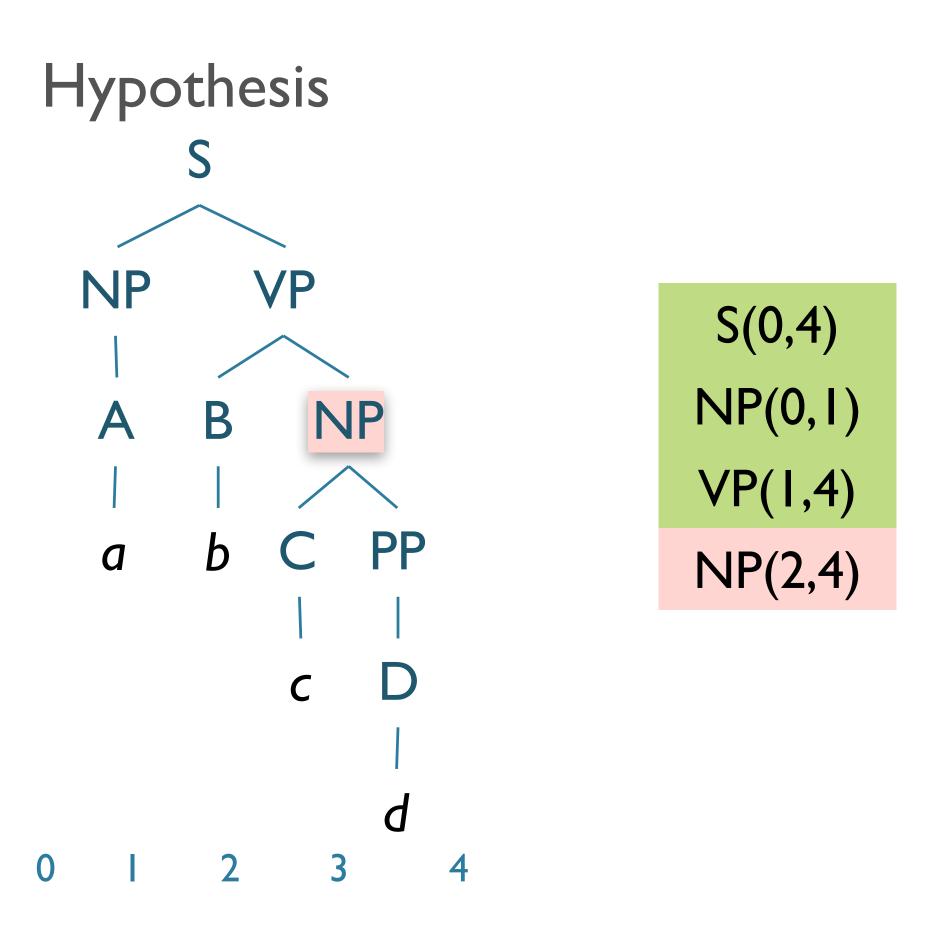


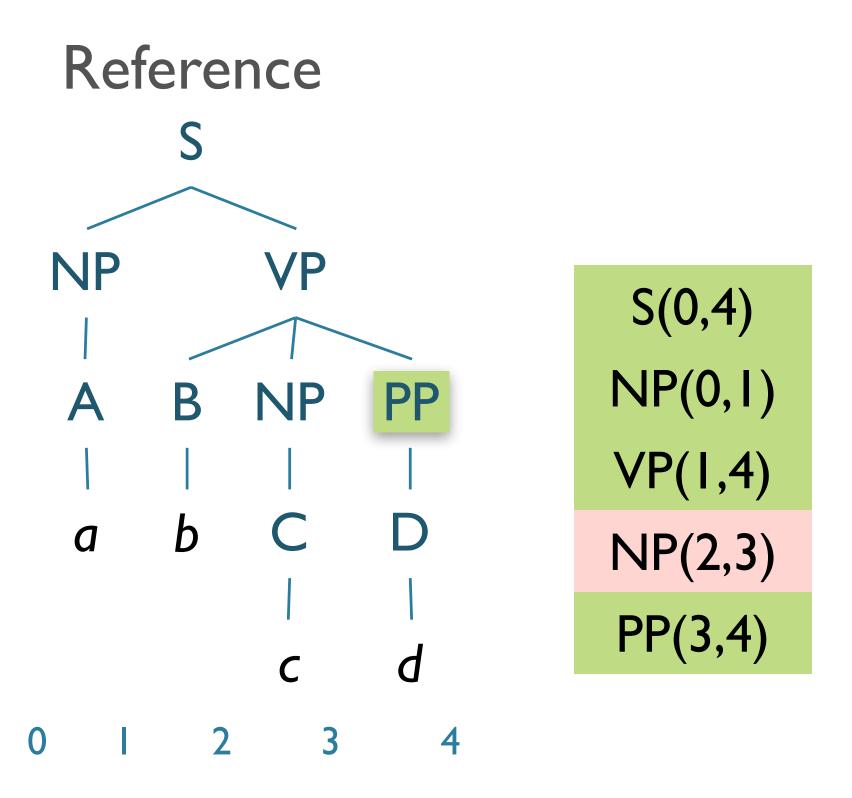


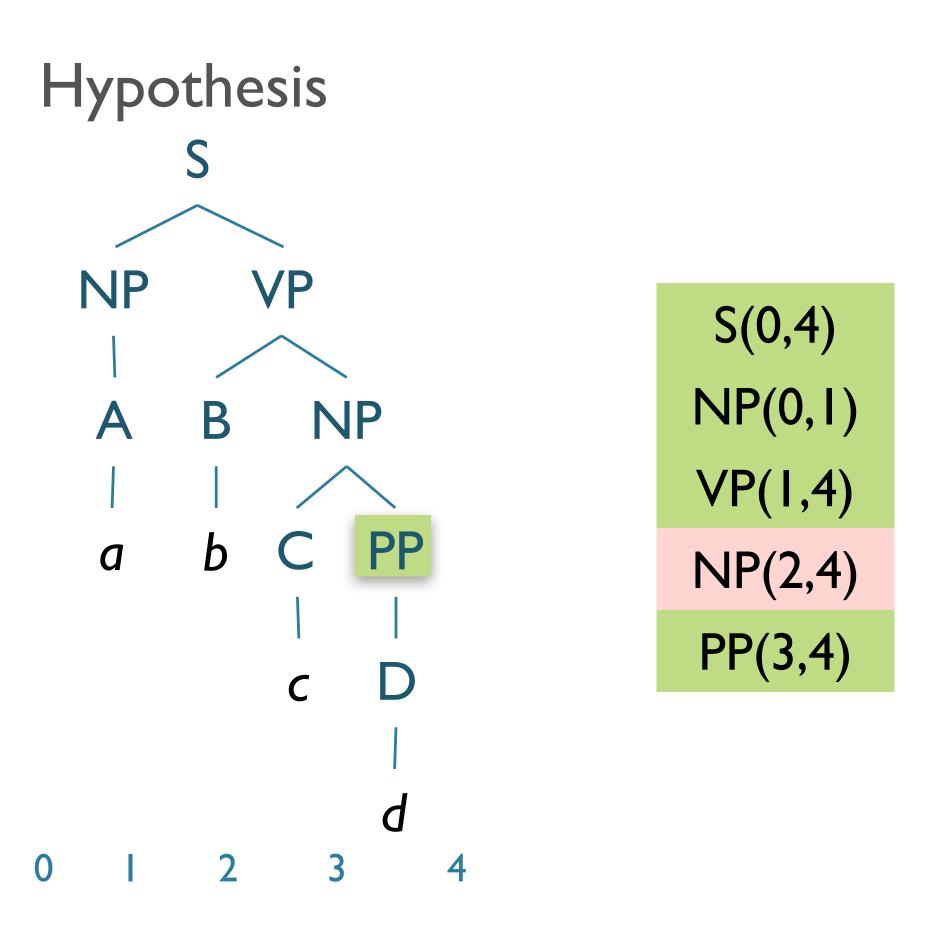




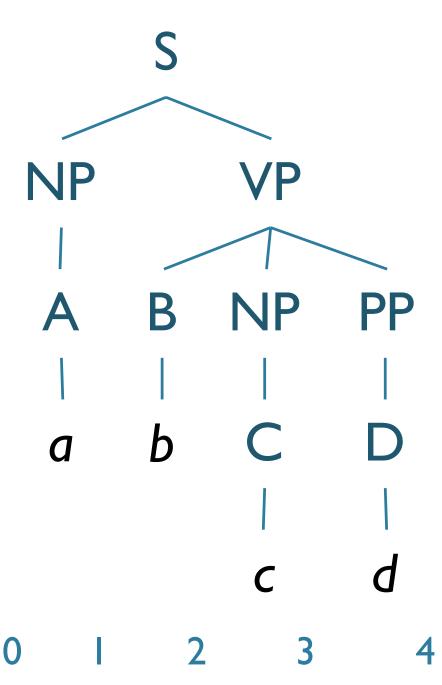








#### Reference



S(0,4) NP(0,I)**VP(1,4)** NP(2,3)PP(3,4)

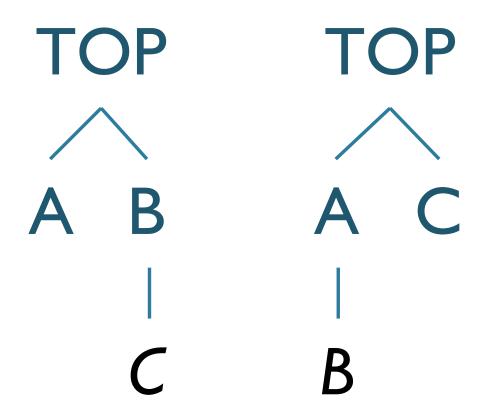
LP:

# Hypothesis **VP** NP NP a

S(0,4) NP(0,1)**VP(1,4)** NP(2,4) PP(3,4)

#### Parser Evaluation

- Crossing Brackets:
  - # of constituents where produced parse has bracketings that overlap for the siblings:
  - ((A B) C) { (0,2), (2,3) } and hyp. has  $(A (B C)) - \{ (0,1), (1,3) \}$



```
/* crossing is counted based on the brackets */
/* in test rather than gold file (by Mike) */
for(j=0;j<bn2;j++){</pre>
 for(i=0;i<bn1;i++){</pre>
    if(bracket1[i].result != 5 &&
       bracket2[j].result != 5 &&
       ((bracket1[i].start < bracket2[j].start &&</pre>
         bracket1[i].end > bracket2[j].start &&
         bracket1[i].end < bracket2[j].end) ||</pre>
        (bracket1[i].start > bracket2[j].start &&
         bracket1[i].start < bracket2[j].end &&</pre>
         bracket1[i].end > bracket2[j].end))){
                 from evalb.c
```

#### State-of-the-Art Parsing

- Parsers trained/tested on Wall Street Journal PTB
  - LR: 94%+;
  - LP: 94%+;
  - Crossing brackets: 1%

- Standard implementation of Parseval:
  - evalb

#### Evaluation Issues

- Only evaluating constituency
- There are other grammar formalisms:
  - LFG (Constraint-based)
  - Dependency Structure
- Extrinsic evaluation
  - How well does getting the correct parse match the semantics, etc?

## Earley Parsing

#### Earley vs. CKY

- CKY doesn't capture full original structure
  - Can back-convert binarization, terminal conversion
  - Unit non-terminals require change in CKY

#### Earley vs. CKY

- CKY doesn't capture full original structure
  - Can back-convert binarization, terminal conversion
  - Unit non-terminals require change in CKY
- Earley algorithm
  - Supports parsing efficiently with arbitrary grammars
  - Top-down search
  - Dynamic programming
    - Tabulated partial solutions
  - Some bottom-up constraints

### Earley Algorithm

- Another dynamic programming solution
  - Partial parses stored in "chart"
  - Compactly encodes ambiguity
  - O(N<sup>3</sup>)
- Chart entries contain:
  - Subtree for a single grammar rule
  - Progress in completing subtree
  - Position of subtree w.r.t. input

#### Earley Algorithm

- First, left-to-right pass fills out a chart with N+1 states
  - Chart entries sit between words in the input string
  - Keep track of states of the parse at those positions
  - For each word position, chart contains set of states representing all partial parse trees generated so far
    - e.g. chart[0] contains all partial parse trees generated at the beginning of sentence

#### **Chart Entries**

- Three types of constituents:
  - Predicted constituents
  - In-progress constituents
  - Completed constituents

#### Parse Progress

- Represented by Dotted Rules
  - Position of indicates type of constituent
- 0 Book 1 that 2 flight 3
  - $\bullet$   $S \rightarrow \bullet VP$ [0,0] (predicted)
  - $NP \rightarrow Det \cdot Nom$  [1,2] (in progress)
  - $VP \rightarrow VNP$  [0,3] (completed)
- [x,y] tells us what portion of the input is spanned so far by rule
- Each state *s<sub>i</sub>*: <*dotted rule*>, [<*back pointer*>, <*current position*>]

#### o Book 1 that 2 flight 3

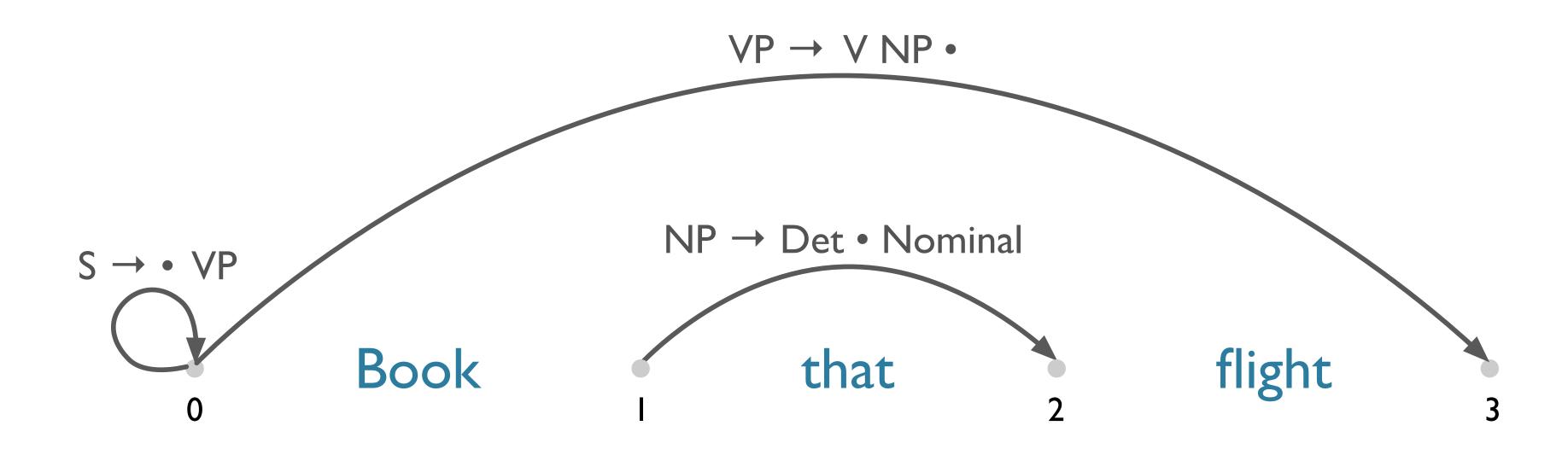
- $S \rightarrow VP$ , [0,0]
  - First 0 means S constituent begins at the start of input
  - Second 0 means the dot is here too
  - So, this is a top-down prediction

#### o Book 1 that 2 flight 3

- $S \to VP, [0,0]$ 
  - First 0 means S constituent begins at the start of input
  - Second 0 means the dot is here too
  - So, this is a top-down prediction
- *NP* → *Det Nom*, [1,2]
  - the NP begins at position 1
  - the dot is at position 2
  - so, Det has been successfully parsed
  - Nom predicted next

#### o Book 1 that 2 flight 3 (continued)

- $VP \rightarrow VNP \cdot [0,3]$ 
  - Successful VP parse of entire input



#### Successful Parse

- Final answer found by looking at last entry in chart
- If entry resembles  $S \rightarrow \alpha \cdot [0,N]$  then input parsed successfully
- Chart will also contain record of all possible parses of input string, given the grammar

### Parsing Procedure for the Earley Algorithm

- Move through each set of states in order, applying one of three operations:
  - predictor: add predictions to the chart
  - scanner: read input and add corresponding state to chart
  - completer: move dot to right when new constituent found
- Results (new states) added to current or next set of states in chart
- No backtracking and no states removed: keep complete history of parse

#### Earley Algorithm

```
function Earley-Parse(words, grammar) returns chart
  ENQUEUE((y \rightarrow \bullet S, [0,0]), chart[0])
 for i \leftarrow from 0 to LENGTH(words) do
   for each state in chart[i] do
      if Incomplete?(state) and
         NEXT-CAT(state) is not a part of speech then
        Predictor(state)
     elseif Incomplete?(state) and
         NEXT-CAT(state) is a part of speech then
        SCANNER(state)
     else
        COMPLETER(state)
     end
    end
  return(chart)
```

#### Earley Algorithm

```
procedure Predictor((A \rightarrow \alpha \cdot B \beta, [i,i]))
  for each (B \rightarrow y) in Grammar-Rules-For(B,grammar) do
     ENQUEUE((B \rightarrow \bullet y, [j,j]), chart[j])
  end
procedure Scanner((A \rightarrow \alpha \cdot B \beta, [i, j]))
  if B ⊂ Parts-of-Speech(word[j]) then
     ENQUEUE((B \rightarrow word[j] \bullet, [j,j+1]), chart[j+1])
procedure Completer((B \rightarrow y \cdot, [j,k]))
 for each (A \rightarrow \alpha \cdot B \beta, [i,j]) in chart[j] do
     ENQUEUE((A \rightarrow \alpha B \cdot \beta, [i,k]), chart[k])
  end
```

#### 3 Main Subroutines of Earley

- Predictor
  - Adds predictions into the chart
- Scanner
  - Reads the input words and enters states representing those words into the chart
- Completer
  - Moves the dot to the right when new constituents are found

#### Predictor

- Intuition:
  - Create new state for top-down prediction of new phrase
- Applied when non part-of-speech non-terminals are to the right of a dot:
  - $S \rightarrow VP[0,0]$
- Adds new states to current chart
  - One new state for each expansion of the non-terminal in the grammar

$$VP \rightarrow \cdot V$$
 [0,0]  
 $VP \rightarrow \cdot V NP$  [0,0]

### Chart[O]

S0	$\gamma \rightarrow \cdot S$	[0,0]	Dummy start state
S1	$S \rightarrow \cdot NP VP$	[0.0]	Predictor
S2	$S \rightarrow \cdot Aux NP VP$	[0,0]	Predictor
S3	$S \rightarrow \cdot VP$	[0,0]	Predictor
S4	NP → • Pronoun NP → • Proper-Noun NP → • Det Nominal	[0,0]	Predictor
S5		[0,0]	Predictor
S6		[0,0]	Predictor
S7 S8 S9 S10 S11	$VP \rightarrow \cdot Verb$ $VP \rightarrow \cdot Verb NP$ $VP \rightarrow \cdot Verb NP PP$ $VP \rightarrow \cdot Verb PP$ $VP \rightarrow \cdot VP PP$	[0,0] [0,0] [0,0] [0,0]	Predictor Predictor Predictor Predictor Predictor

### Chart[1]

S12	Verb → book •	[0,1]	Scanner
S13 S14 S15 S16	$VP \rightarrow Verb \cdot NP$ $VP \rightarrow Verb \cdot NP$ $VP \rightarrow Verb \cdot NP PP$ $VP \rightarrow Verb \cdot PP$	[0,1] [0,1] [0,1]	Completer Completer Completer Completer
S17	$S \rightarrow VP$ .	[0,1]	Completer
S18	$VP \rightarrow VP \cdot PP$	[0,1]	Completer
S19 S20 S21 S22	$NP \rightarrow \cdot Pronoun$ $NP \rightarrow \cdot Proper-Noun$ $NP \rightarrow \cdot Det Nominal$ $PP \rightarrow \cdot Prep NP$	[1,1] [1,1] [1,1] [1,1]	Predictor Predictor Predictor Predictor

S0: 
$$y \rightarrow S[0,0]$$



S0: 
$$y \rightarrow S[0,0]$$

S3: 
$$S \rightarrow VP[0,0]$$



```
S0: y \rightarrow S[0,0]
```

S3:  $S \rightarrow VP[0,0]$ 

S8:  $VP \rightarrow \bullet Verb NP [0,0]$ 

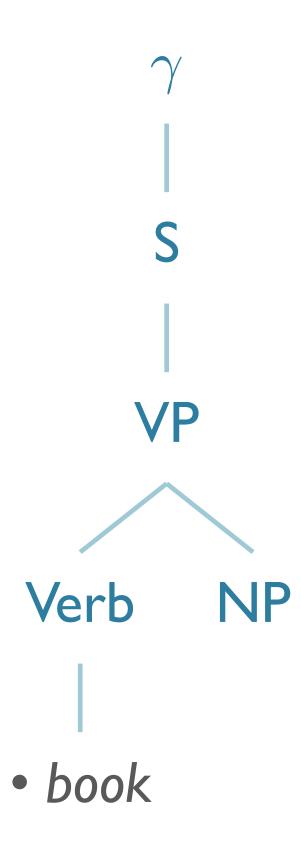


```
S0: y \rightarrow S[0,0]
```

S3:  $S \rightarrow VP[0,0]$ 

S8:  $VP \rightarrow \cdot Verb NP [0,0]$ 

S12:  $Verb \rightarrow \cdot book [0,0]$ 

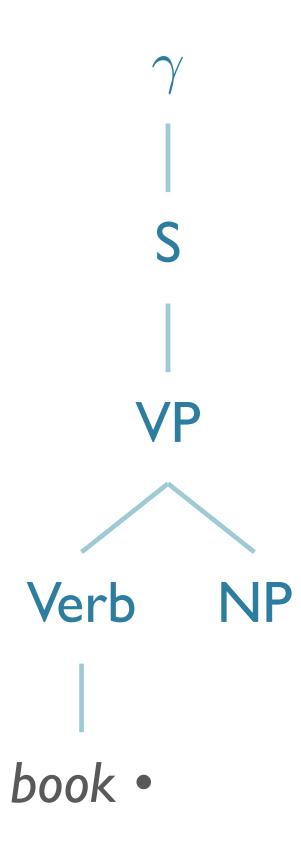


```
S0: y \rightarrow S[0,0]
```

S3:  $S \rightarrow VP[0,0]$ 

S8:  $VP \rightarrow \cdot Verb NP [0,0]$ 

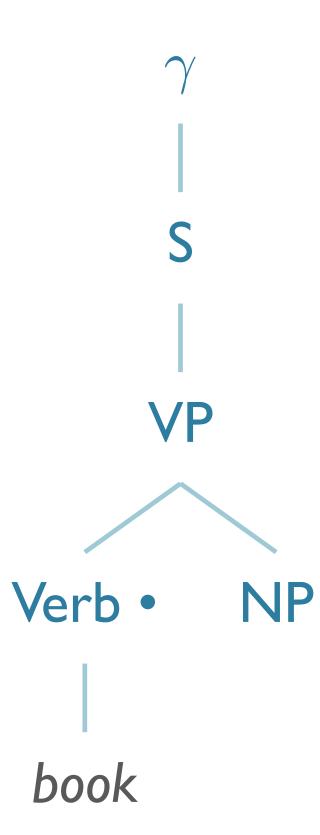
S12:  $Verb \rightarrow book \cdot [0,1]$ 



```
S0: y \rightarrow S[0,0]
```

S3:  $S \rightarrow VP[0,0]$ 

S8:  $VP \rightarrow Verb \cdot NP$  [0,1]



```
S0: y \rightarrow S[0,0]
```

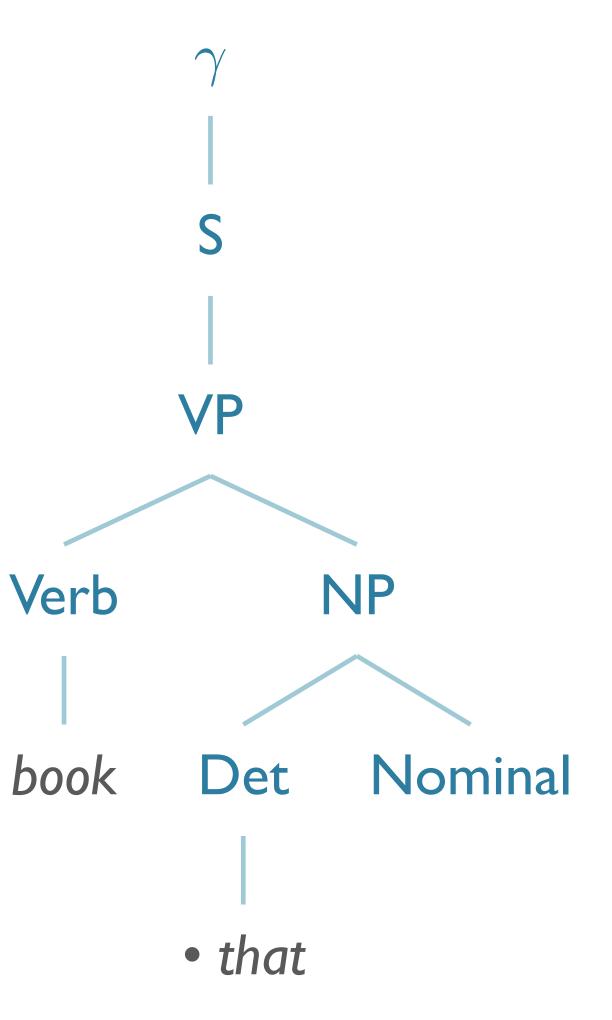
S3:  $S \rightarrow VP \cdot [0,1]$ 

S8:  $VP \rightarrow Verb \cdot NP$  [0,1]

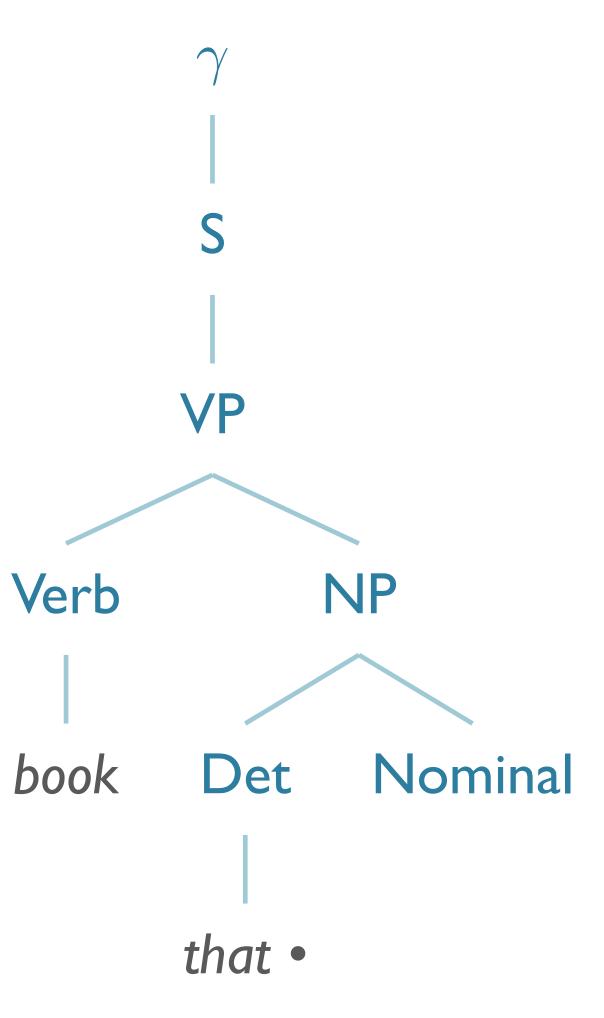


```
S0: y \rightarrow S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP \rightarrow Verb \cdot NP [0,1]
S21: NP \rightarrow \bullet Det Nominal [1,1]
                                                     VP
                                            Verb
                                                             NP
                                            book
                                                     Det
                                                                Nominal
```

```
S0: y \rightarrow S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP \rightarrow Verb \cdot NP [0,1]
S21: NP \rightarrow \bullet Det Nominal [1,1]
S23: Det → • "that" [1,1]
```



```
S0: y \rightarrow S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP \rightarrow Verb \cdot NP [0,1]
S21: NP \rightarrow \bullet Det Nominal [1,1]
S23: Det → "that" • [1,2]
```



```
S0: y \rightarrow S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP \rightarrow Verb \cdot NP [0,1]
S21: NP \rightarrow Det \cdot Nominal [1,2]
                                                     VP
                                            Verb
                                                             NP
                                                                Nominal
                                            book
                                                      Det •
                                                       that
```

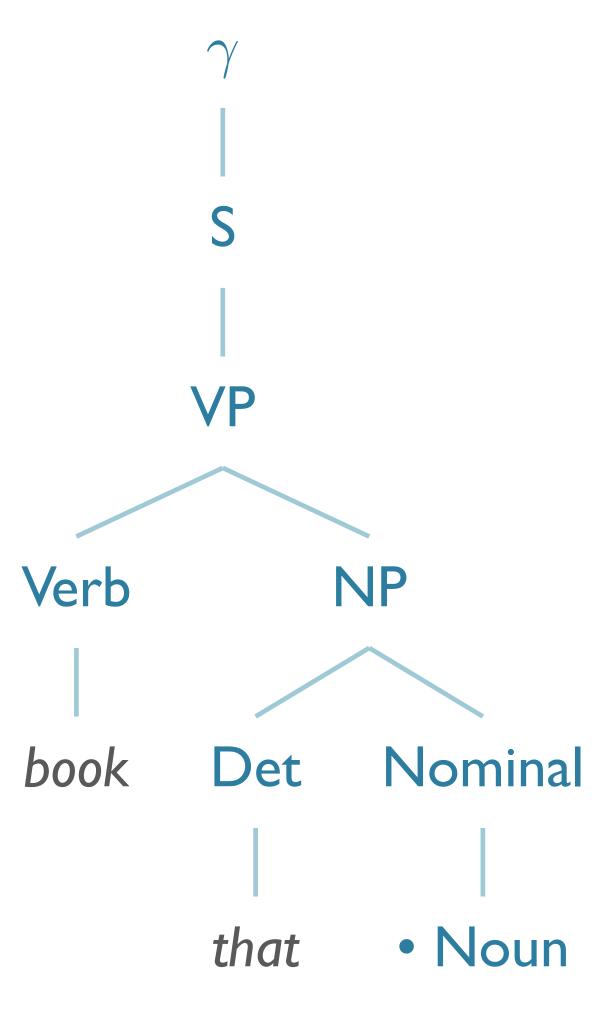
```
S0: y \rightarrow S[0,0]
```

S3:  $S \rightarrow VP \cdot [0,1]$ 

S8:  $VP \rightarrow Verb \cdot NP$  [0,1]

S21:  $NP \rightarrow Det \cdot Nominal [1,2]$ 

S25: Nominal  $\rightarrow$  • Noun [2,2]



```
S0: y \rightarrow S[0,0]
```

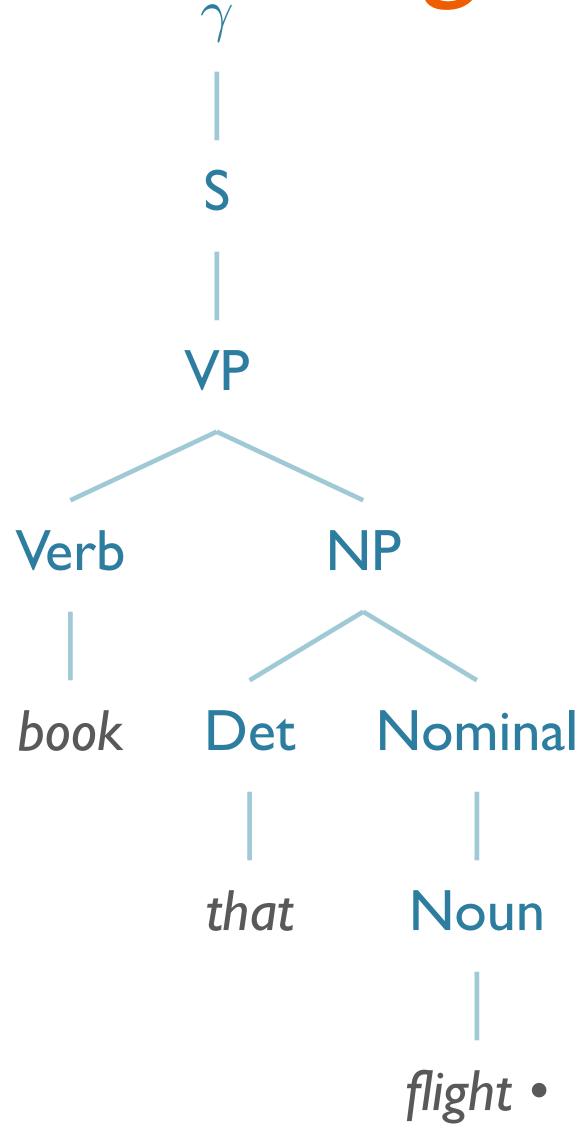
S3: 
$$S \rightarrow VP \cdot [0,1]$$

S8: 
$$VP \rightarrow Verb \cdot NP [0,1]$$

S21:  $NP \rightarrow Det \cdot Nominal$  [1,2]

S25: *Nominal* → • *Noun* [2,2]

S28: *Noun* → "*flight*" • [2,3]



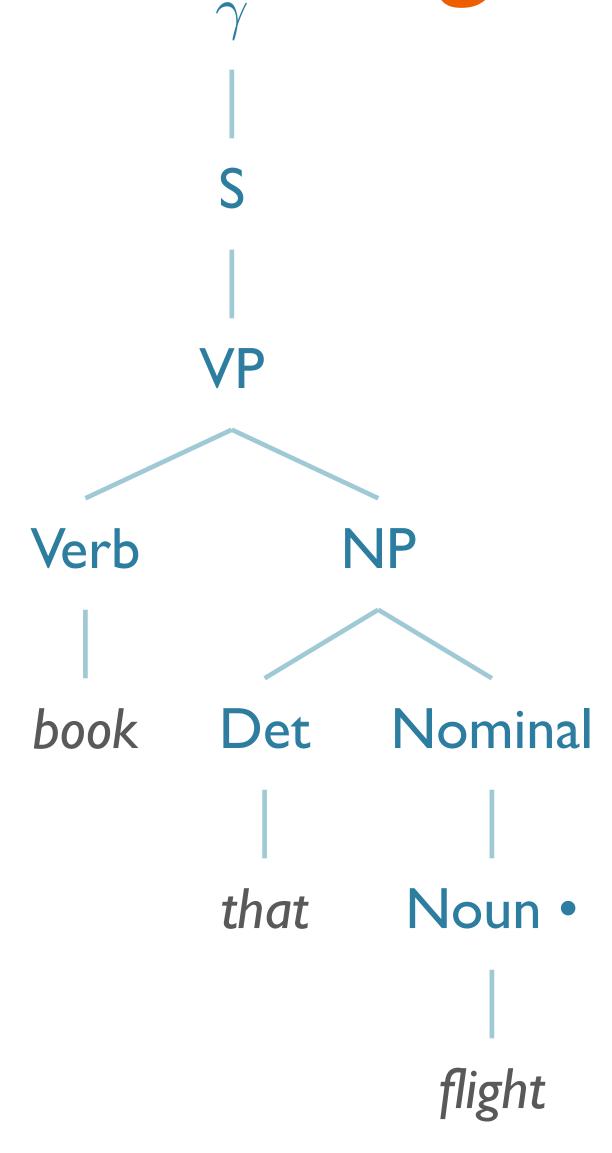
```
S0: y \rightarrow S[0,0]
```

S3: 
$$S \rightarrow VP \cdot [0,1]$$

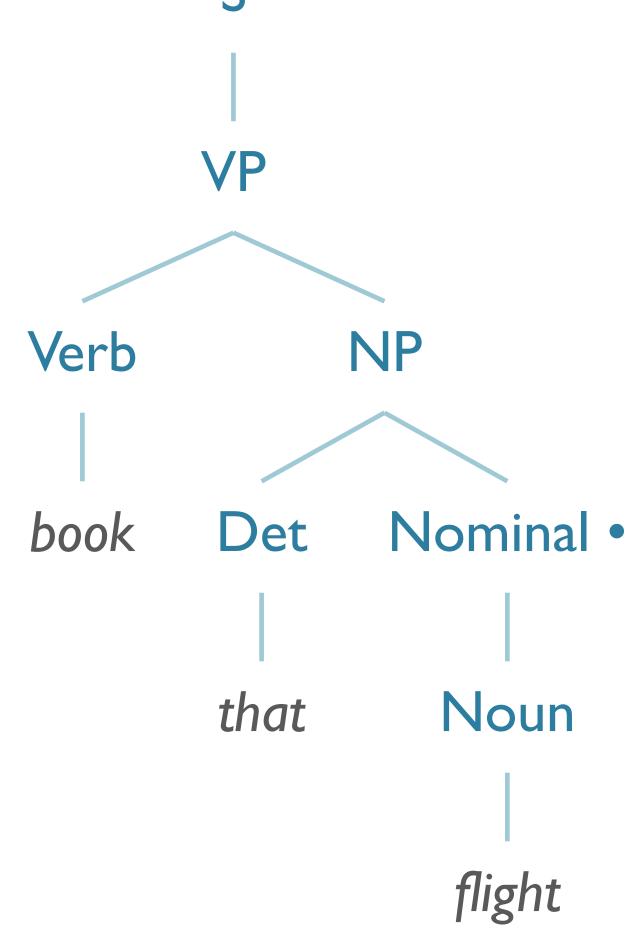
S8: 
$$VP \rightarrow Verb \cdot NP$$
 [0,1]

S21:  $NP \rightarrow Det \cdot Nominal$  [1,2]

S25: Nominal  $\rightarrow$  Noun • [2,3]



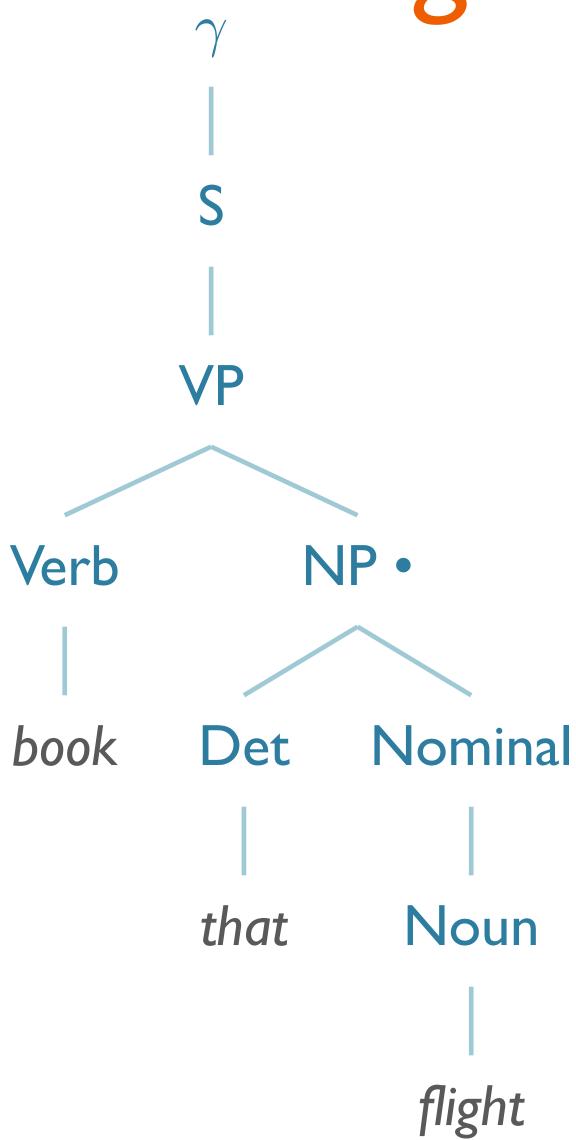
S0:  $y \rightarrow \cdot S$  [0,0] S3:  $S \rightarrow VP \cdot [0,1]$ S8:  $VP \rightarrow Verb \cdot NP$  [0,1] S21:  $NP \rightarrow Det Nominal \cdot [1,3]$ 



S0:  $y \rightarrow S[0,0]$ 

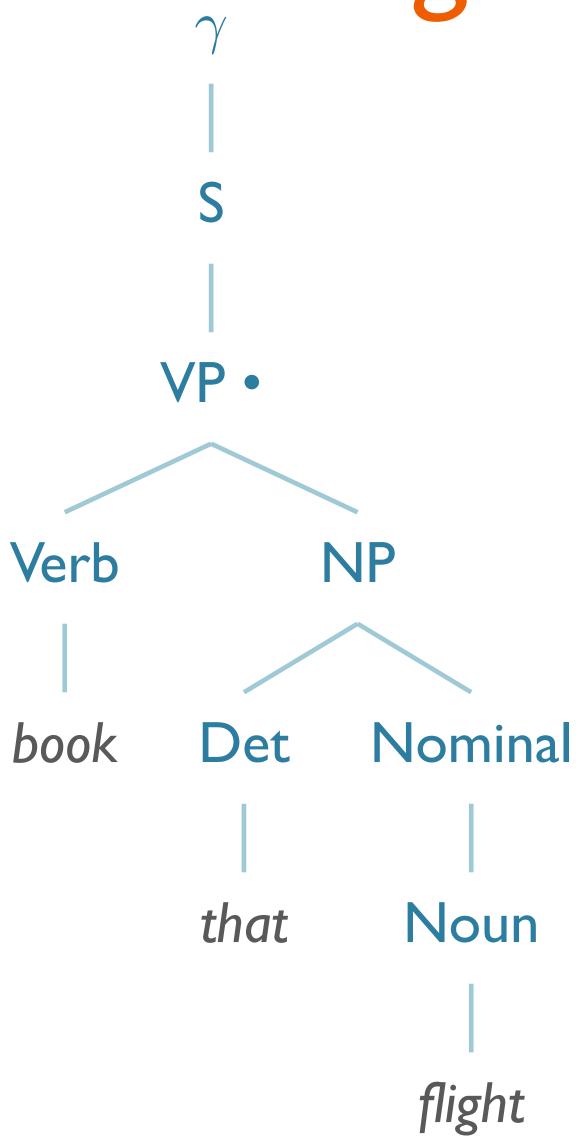
S3:  $S \rightarrow VP \cdot [0,1]$ 

S8:  $VP \rightarrow Verb NP \cdot [0,3]$ 



S0:  $y \rightarrow S[0,0]$ 

S3:  $S \rightarrow VP \cdot [0,3]$ 



### What About Dead Ends?

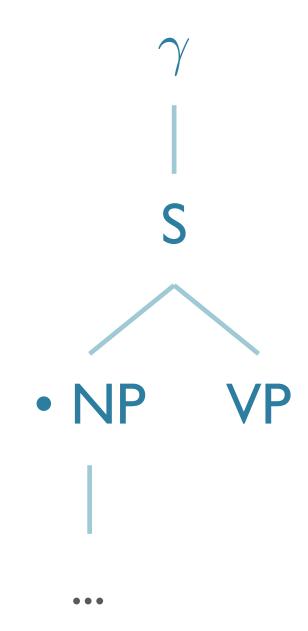
S0:  $y \rightarrow S[0,0]$ 

 $S1: S \rightarrow PVP[0,0]$ 

NP → • Pronoun

*NP* → • *Proper-Noun* 

NP → • Det Nominal

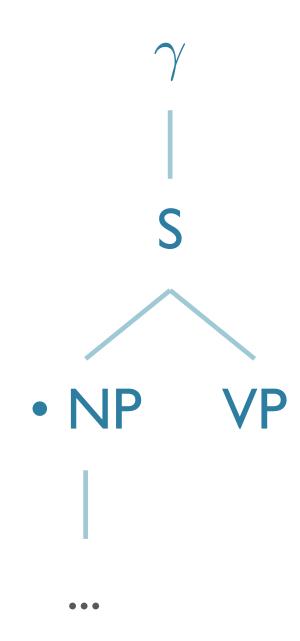


book

S0: 
$$y \rightarrow S[0,0]$$

 $S1: S \rightarrow PVP[0,0]$ 

$$NP \rightarrow Pronoun$$



book

 We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?

- We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?
- No!

```
procedure Enqueue(state, chart-entry)
 if state is not already in chart-entry then
    Push(state, chart-entry)
  end
```

- We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?
- No!

```
procedure Enqueue(state, chart-entry)
 if state is not already in chart-entry then
    Push(state, chart-entry)
  end
```

**Exercise**: parse 'table and chair' using the very simple grammar Nom -> Nom 'and' Nom | 'table' | 'chair'