PCFGs: Parsing & Evaluation

LING 571 — Deep Processing Techniques for NLP Shane Steinert-Threlkeld







Roadmap

- CKY + back-pointers
- PCFGs
- PCFG Parsing (PCKY)
- Inducing a PCFG
- Evaluation
- [Earley parsing]
- HW3 + collaboration







CKY Parsing: Backpointers





Current CKY Algorithm

Limitations:

Only stores non-terminals in cell Not rules or cells corresponding to RHS Stores SETS of non-terminals Multiple rules with same LHS collide

Currently only *acceptance/recognition*





Backpointers

- Instead of list of possible nonterminals for that node, each cell should have:
 - Nonterminal for the node
 - Pointer to left and right children cells
 - Either direct pointer to cell, or indices

- example:
- ackPointer()
- d = [X2, (1,4)]
- d = [PP, (4, 6)]









• Pair each nonterminal with back-pointer to cells from which it was derived

• Last step:

• construct trees from back-pointers in [0, n]

CKY Parser







NP, Pronoun [0,1]

S		S		
[0,2]	[0,3]	[0,4]	[0,5]	
Verb, VP, S		VP, X2, S		VP
[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
	Det	NP		NP
	[2,3]	[2,4]	[2,5]	[2,6]
		Noun, Nom		Nom
		[3,4]	[3,5]	[3,6]
			Prep	PP
			[4,5]	[4,6]
				NNP, NP
				[5,6]







bp_1 = BackPointer() $bp_{1.l}child = [VP, (1,4)]$ $bp_1.r_child = [PP, (4, 6)]$

S		S		
[0,2]	[0,3]	[0,4]	[0,5]	
Verb, VP, S		VP, X2, S		VP
[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
	Det	NP		NP
	[2,3]	[2,4]	[2,5]	[2,6]
		Noun, Nom		Nom
		[3,4]	[3,5]	[3,6]
			Prep	PP
			[4,5]	[4,6]
				NNP, NP
				[5,6]







bp_2 = BackPointer() $bp_2.l_child = [X2, (1,4)]$ $bp_2.r_child = [PP, (4,6)]$

S		S		
[0,2]	[0,3]	[0,4]	[0,5]	
Verb, VP, S		VP, X2, S		VP
[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
	Det	NP		NP
	[2,3]	[2,4]	[2,5]	[2,6]
		Noun, Nom		Nom
		[3,4]	[3,5]	[3,6]
			Prep	PP
			[4,5]	[4,6]
				NNP, NP
				[5,6]













CKY Discussion

- Running time:

 - $O(n^3)$ where n is the length of the input string • Inner loop grows as square of # of non-terminals
- Expressiveness:
 - As implemented, requires CNF Weak equivalence to original grammar

 - Doesn't capture full original structure
 - Back-conversion?
 - Can do binarization, terminal conversion • Unit productions requires change in CKY





CKY + Back-pointers Example









l prefer a

	S		S		S
	[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
	Verb, VP, S		VP, X2, S		VP, X2, S
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
	S	[2,3]	[2,4]	[2,5]	[2,6]
	\sim		Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
INF	V F			Prep	PP
				[4,5]	[4,6]
					NNP, NP
					[5,6]









prefer a

	S		S		S
	[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
	Verb, VP, S		VP, X2, S		VP, X2, S
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
	S	[2,3]	[2,4]	[2,5]	[2,6]
	\sim		Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
				Prep	PP
				[4,5]	[4,6]
1					NNP, NP
					[5,6]





cky_table[0,6][S] NP, (0,1),Pronoun **VP**, (1,6)) [0,1] cky_table [,1][NP] $= \{ (' I ') \}$ cky_table[<mark>1,6</mark>][VP] {(Verb, (1,2), = NP, (2,6)), (X2, (1,4), PP, (4,6))}

þrefer a

	S		S		S
un	[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
	Verb, VP, S		VP, X2, S		VP, X2, S
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
	S	[2,3]	[2,4]	[2,5]	[2,6]
			Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
	V F			Prep	PP
				[4,5]	[4,6]
1	verb NP				NNP, NP
					[5,6]

flight

TWA







NP, $cky_table[0,6][S] = {(NP, (0,1),$ Pronoun VP, (1,6)) [0,1] $cky_table[0,1][NP] = {('I')}$ $cky_table[1,6][VP] = {(Verb, (1,2),$ NP, (2,6)), (X2, (1,4),PP, (4,6))} $cky_table[1,2][Verb] = {('prefer')}$

prefer a

	S		S		S
un	[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
	Verb, VP, S		VP, X2, S		VP, X2, S
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
	S	[2,3]	[2,4]	[2,5]	[2,6]
			Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
	۷۲			Prep	PP
				[4,5]	[4,6]
Ι	Verb NP				NNP, NP
					[5,6]
	þrefer				

TWA







NP, $cky_table[0,6][S] = {(NP, (0,1),$ Pronoun VP, (1,6)) [0,1] $cky_table[0,1][NP] = {('I')}$ $cky_table[1,6][VP] = {(Verb, (1,2),$ NP, (2,6)), (X2, (1,4),PP, (4,6))} cky_table[1,2][Verb] = {('prefer')} $cky_table[2,6][NP] = {(Det, (2,3),$ NP Nom, (3,6)}

prefer a

S		S		S
[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
Verb, VP, S		VP, X2, S		VP, X2, S
[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
	Det	NP		NP
S	[2,3]	[2,4]	[2,5]	[2,6]
		Noun, Nom		Nom
		[3,4]	[3,5]	[3,6]
VP		[3,4]	[3,5] Prep	[3,6] PP
VP		[3,4]	[3,5] Prep [4,5]	[3,6] PP [4,6]
VP Verb	NP	[3,4]	[3,5] Prep [4,5]	[3,6] PP [4,6] NNP, NP
VP Verb	NP	[3,4]	[3,5] Prep [4,5]	[3,6] PP [4,6] NNP, NP [5,6]

flight TWA on









NP, $cky_table[0,6][S] = {(NP, (0,1),$ Pronoun VP, (1,6)) [0,1] $cky_table[0,1][NP] = {('I')}$ $cky_table[1,6][VP] = {(Verb, (1,2),$ NP, (2,6)), (X2, (1,4),PP, (4,6))} cky_table[1,2][Verb] = {('prefer')} $cky_table[2,6][NP] = {(Det, (2,3),$ NP Nom, (3, 6) $cky_table[2,3][Det] = {('a')}$

	S		S		S
	[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
	Verb, VP, S		VP, X2, S		VP, X2, S
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
	S	[2,3]	[2,4]	[2,5]	[2,6]
			Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
	VP			Prep	PP
_				[4,5]	[4,6]
	Verb	NP			NNP, NP
					[5,6]
Þ	orefer Det	Nom			
	а				

TWA

flight on









NP, $cky_table[0,6][S] = {(NP, (0,1),$ Pronoun VP, (1,6)) [0,1] $cky_table[0,1][NP] = {('I')}$ $cky_table[1,6][VP] = {(Verb, (1,2),$ NP, (2,6)), (X2, (1,4),PP, (4,6))} cky_table[1,2][Verb] = {('prefer')} $cky_table[2,6][NP] = {(Det, (2,3),$ NP Nom, (3, 6) $cky_table[2,3][Det] = {('a')}$

	S		S		S
	[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
	Verb, VP, S		VP, X2, S		VP, X2, S
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
	S	[2,3]	[2,4]	[2,5]	[2,6]
			Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
	VP			Prep	PP
_				[4,5]	[4,6]
	Verb	NP			NNP, NP
					[5,6]
Þ	orefer Det	Nom			
	а				

TWA

flight on









a

	S		S		S
	[0,2]	[0,3]	[0,4]	[0,5]	[0,6]
	Verb, VP, S		VP, X2, S		VP, X2, S
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
		[2,3]	[2,4]	[2,5]	[2,6]
			Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
	C			Prep	PP
	2			[4,5]	[4,6]
					NNP, NP
NP	VP				[5,6]

flight

1

X2 PP

TWA







Probabilistic Context-Free Grammars





Probabilistic Context-free Grammars: Roadmap

Motivation: Ambiguity

Approach:

Definition

Disambiguation

Parsing

Evaluation

Enhancements







Motivation

What about ambiguity?

Current algorithm can *represent* it

...can't resolve it.

 \mathbf{W} university of washington





Probabilistic Parsing

- Provides strategy for solving disambiguation problem
 - Compute the probability of all analyses • Select the most probable

- Employed in language modeling for speech recognition
 - N-gram grammars predict words, constrain search
 - Also, constrain generation, translation





N

a set of **non-terminal symbols** (or **variables**)









a set of **non-terminal symbols** (or **variables**)

a set of **terminal symbols** (disjoint from N)







a set of non-ter	N
a set of termi	\sum
a set of rules of productions, each of t A is a non-terminal, $meta$ is a string of sy is a number betw	R

minal symbols (or **variables**)

nal symbols (disjoint from N)

the form $A \rightarrow \beta[p]$, where A is a non-terminal where ymbols from the infinite set of strings $(\Sigma \cup N)*$ and pween 0 and 1 expressing P(eta|A)







a set of non-ter	N
a set of termi	\sum
a set of rules of productions, each of t A is a non-terminal, eta is a string of sy	R
is a number betw	
a desig	S

minal symbols (or **variables**)

nal symbols (disjoint from N)

the form $A \rightarrow \beta[p]$, where A is a non-terminal where

ymbols from the infinite set of strings $(\Sigma \cup N)*$ and p

ween 0 and 1 expressing P(eta|A)

a designated start symbol







- will be expanded as RHS
 - $P(A \rightarrow \beta)$
 - $P(A \rightarrow \beta | A)$
 - $P(\boldsymbol{\beta}|A)$
 - $P(RHS \mid LHS)$
- really meant.

PCFGs

Augment each production with probability that LHS

• NB: the first is often used; but the latter are what's







• Sum over all possible expansions is 1



- is 1
 - Recursive rules often yield inconsistent grammars (Booth & Thompson, 1973)

PCFGs

$\sum_{\alpha} P(A \to \beta) = 1$

• A PCFG is *consistent* if sum of probabilities of all sentences in language







Example PCFG: Augmented \mathcal{L}_1

Grammar	
$S \rightarrow NP VP$	[.80]
$S \rightarrow Aux NP VP$	[.15]
$S \rightarrow VP$	[.05]
$NP \rightarrow Pronoun$	[.35]
$NP \rightarrow Proper-Noun$	[.30]
$NP \rightarrow Det Nominal$	[.20]
$NP \rightarrow Nominal$	[.15]
$Nominal \rightarrow Noun$	[.75]
$Nominal \rightarrow Nominal Noun$	[.20]
$Nominal \rightarrow Nominal PP$	[.05]
$VP \rightarrow Verb$	[.35]
$VP \rightarrow Verb NP$	[.20]
$VP \rightarrow Verb \ NP \ PP$	[.10]
$VP \rightarrow Verb PP$	[.15]
$VP \rightarrow Verb NP NP$	[.05]
$VP \rightarrow VP PP$	[.15]
$PP \rightarrow Preposition NP$	[1.0]

Lexicon

 $Det \rightarrow that [.10] \mid a [.30] \mid the [.60]$ Noun \rightarrow book [.10] | flight [.30] | meal [.15] | money [0.5] | *flights* [0.40] | *dinner* [.10] $Verb \rightarrow book [.30] \mid include [.30] \mid prefer [.40]$ $Pronoun \rightarrow I[.40] \mid she \mid .05] \mid me \mid .15] \mid you \mid .40]$ Proper-Noun \rightarrow Houston [.60] | NWA [.40] $Aux \rightarrow does [.60] \mid can [.40]$ $Preposition \rightarrow from [.30] \mid to [.30] \mid on [.20] \mid near [.15]$ through [.05]









Example PCFG: Augmented \mathcal{L}_1

Grammar	
$S \rightarrow NP VP$	[.80]
$S \rightarrow Aux NP VP$	[.15]
$S \rightarrow VP$	[.05]
$NP \rightarrow Pronoun$	[.35]
$NP \rightarrow Proper-Noun$	[.30]
$NP \rightarrow Det Nominal$	[.20]
$NP \rightarrow Nominal$	[.15]
$Nominal \rightarrow Noun$	[.75]
$Nominal \rightarrow Nominal Noun$	[.20]
$Nominal \rightarrow Nominal PP$	[.05]
$VP \rightarrow Verb$	[.35]
$VP \rightarrow Verb NP$	[.20]
$VP \rightarrow Verb \ NP \ PP$	[.10]
$VP \rightarrow Verb PP$	[.15]
$VP \rightarrow Verb NP NP$	[.05]
$VP \rightarrow VP PP$	[.15]
$PP \rightarrow Preposition NP$	[1.0]

Lexicon

 $Det \rightarrow that [.10] \mid a [.30] \mid the [.60]$ Noun \rightarrow book [.10] | flight [.30] | meal [.15] | money [0.5] | *flights* [0.40] | *dinner* [.10] $Verb \rightarrow book [.30] \mid include [.30] \mid prefer [.40]$ $Pronoun \rightarrow I[.40] \mid she \mid .05] \mid me \mid .15] \mid you \mid .40]$ Proper-Noun \rightarrow Houston [.60] | NWA [.40] $Aux \rightarrow does [.60] \mid can [.40]$ $Preposition \rightarrow from [.30] \mid to [.30] \mid on [.20] \mid near [.15]$ through [.05]







Disambiguation

- A PCFG assigns probability to each parse tree T for input S
- Probability of T: product of all rules used to derive T



 $P(T,S) = \prod P(RHS_i | LHS_i)$ $P(T, S) = P(T)P(S \mid T) = P(T)$





	S				
NP			VP		
Pron	Verb	NP			PP
	þrefer	Det a	Nom Noun flight	P on	NP NNP TWA
		S → N NP → Pron	NPVP Pron → I		[0.8] [0.35] [0.4]
	1	$VP \rightarrow V NP PP$ $V \rightarrow prefer$ $NP \rightarrow Det Nom$ $Det \rightarrow a$ $Nom \rightarrow N$ $N \rightarrow flight$ $PP \rightarrow P NP$ $P \rightarrow on$ $NP \rightarrow NNP$ $NNP \rightarrow NVP$			[0.1] [0.4] [0.2] [0.3] [0.75] [0.3] [1.0] [0.2] [0.3] [0.3] [0.4]







	S				
NP			VP		
Pron	Verb	NP			PP
	prefer	Det	Nom	P	
		а	Noun flight	on	NNP TWA
		S → N NP → Pron	NPVP Pron → I		[0.8] [0.35] [0.4]
		$VP \rightarrow V$	NP PP		[0.1]
		$\wedge \rightarrow \downarrow$	orefer		[0.4]
		$NP \rightarrow Det Nom$			[0.2] [0.3]
		Nom \rightarrow N			[0.75]
		$N \rightarrow$	flight		[0.3]
		PP →	P NP		[1.0]
			on		[0.2]
					[U.3] [0 4]
					נד.ט]

~1.452 × 10⁻⁶









Parsing Problem for PCFGs

• Select T such that (s.t.)



- String of words S is *yield* of parse tree
- Select the tree $\overline{7}$ that maximizes the probability of the parse

 $\hat{T}(S) = \arg \max P(T)$ T s.t. S=yield(T)






• *n*-grams helpful for modeling the probability of a string







- *n*-grams helpful for modeling the probability of a string
- To model a whole sentence with *n-grams* either:







- *n*-grams helpful for modeling the probability of a string
- To model a whole sentence with *n-grams* either:
 - Must use 10+-grams... too sparse







- *n*-grams helpful for modeling the probability of a string
- To model a whole sentence with *n-grams* either:
 - Must use 10+-grams... too sparse
 - Approximate using conditioning on limited context:

$$P(w_i | w_{i-1}) = \frac{P(w_{i-1}, w_i)}{P(w_{i-1})}$$







- *n*-grams helpful for modeling the probability of a string
- To model a whole sentence with *n-grams* either:
 - Must use 10+-grams... too sparse
 - Approximate using conditioning on limited context: $P(w_i | w_{i-1}) = \frac{P(w_{i-1}, w_i)}{P(w_{i-1})}$
- PCFGs are able to give probability of entire string without as bad sparsity





- *n*-grams helpful for modeling the probability of a string
- To model a whole sentence with *n-grams* either:
 - Must use 10+-grams... too sparse
 - Approximate using conditioning on limited context: $P(w_i | w_{i-1}) = \frac{P(w_{i-1}, w_i)}{P(w_{i-1})}$
- PCFGs are able to give probability of entire string without as bad sparsity • Model probability of *syntactically valid* sentences





- *n*-grams helpful for modeling the probability of a string
- To model a whole sentence with *n-grams* either:
 - Must use 10+-grams... too sparse
 - Approximate using conditioning on limited context: $P(w_i | w_{i-1}) = \frac{P(w_{i-1}, w_i)}{P(w_{i-1})}$
- PCFGs are able to give probability of entire string without as bad sparsity
- Model probability of *syntactically valid* sentences
 - Not just probability of sequence of words





PCFGs: Parsing







Probabilistic CKY (PCKY)

- Like regular CKY

•
$$A \rightarrow B C$$

• $A \rightarrow W$

- Represent input with indices b/t words:
 - Book 1 that 2 flight 3 through 4 Houston 5

• Assumes grammar in Chomsky Normal Form (CNF)







Probabilistic CKY (PCKY)

- For input string length n and non-terminals V
 - Cell [i, j, A] in $(n+1) \times (n+1) \times V$ matrix
 - Contains probability that A spans [i, j]







for $j \leftarrow from 1$ to LENGTH(*words*) do for all { $A \mid A \rightarrow words[j] \in grammar$ } $table[j-1, j, A] \leftarrow P(A \rightarrow words[j])$ for $i \leftarrow$ from j-2 downto 0 do for $k \leftarrow i + 1$ to j - 1 do for all $\{A \mid A \rightarrow B \ C \in grammar, \}$ and table[i, k, B] > 0 and table[k, j, C] > 0 } if $(table[i, j, A] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C])$ then $table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]$ $back[i, j, A] \leftarrow \{k, B, C\}$ return BUILD_TREE(back[1, LENGTH(words), S]), table[1,LENGTH(words), S]

PCKY Algorithm









for $j \leftarrow from 1$ to LENGTH(*words*) do for all { $A \mid A \rightarrow words[j] \in grammar$ } $table[j-1, j, A] \leftarrow P(A \rightarrow words[j])$ for $i \leftarrow$ from j-2 downto 0 do for $k \leftarrow i + 1$ to j - 1 do for all $\{A \mid A \rightarrow B \ C \in grammar, \}$ and table[i, k, B] > 0 and table[k, j, C] > 0 } if $(table[i, j, A] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C])$ then $table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]$ $back[i, j, A] \leftarrow \{k, B, C\}$ return BUILD TREE(back 1, LENGTH(words), S]), table 1,LENGTH(words), S]

PCKY Algorithm







for $j \leftarrow from 1$ to LENGTH(*words*) do for all { $A \mid A \rightarrow words[j] \in grammar$ } $table[j-1, j, A] \leftarrow P(A \rightarrow words[j])$ for $i \leftarrow$ from j-2 downto 0 do for $k \leftarrow i + 1$ to j - 1 do for all $\{A \mid A \rightarrow B \ C \in grammar, \}$ and table[i, k, B] > 0 and table[k, j, C] > 0if $(table[i, j, A] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C])$ then $table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]$ $back[i, j, A] \leftarrow \{k, B, C\}$ **return** BUILD TREE(*back*[1, LENGTH(*words*), S]), *table*[1,LENGTH(*words*), S]

PCKY Algorithm







for $j \leftarrow from 1$ to LENGTH(*words*) do for all { $A \mid A \rightarrow words[j] \in grammar$ } $table[j-1, j, A] \leftarrow P(A \rightarrow words[j])$ for $i \leftarrow$ from j-2 downto 0 do for $k \leftarrow i + 1$ to j - 1 do for all $\{A \mid A \rightarrow B \ C \in grammar, \}$ and table[i, k, B] > 0 and table[k, j, C] > 0if $(table[i, j, A] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C])$ then $table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]$ $back[i, j, A] \leftarrow \{k, B, C\}$ **return** BUILD TREE(*back* [1, LENGTH(*words*), S]), *table* [1,LENGTH(*words*), S]

PCKY Algorithm







for $j \leftarrow from 1$ to LENGTH(*words*) do for all { $A \mid A \rightarrow words[j] \in grammar$ } $table[j-1, j, A] \leftarrow P(A \rightarrow words[j])$ for $i \leftarrow$ from j-2 downto 0 do for $k \leftarrow i + 1$ to j - 1 do for all $\{A \mid A \rightarrow B \ C \in grammar, \}$ and table[i, k, B] > 0 and table[k, j, C] > 0 } if $(table[i, j, A] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C])$ then $table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]$ $back[i, j, A] \leftarrow \{k, B, C\}$ **return** BUILD TREE(*back* [1, LENGTH(*words*), S]), *table* [1,LENGTH(*words*), S]

PCKY Algorithm







PCKY Grammar Segment

 $S \rightarrow NP VP \quad [0.80]$ $NP \rightarrow Det N$ [0.30] $VP \rightarrow VNP$ [0.20]

$Det \rightarrow the$	[0.40]
$Det \rightarrow a$	[0.40]
$V \rightarrow includes$	[0.05]
$N \rightarrow meal$	[0.01]
$N \rightarrow flight$	[0.02]







Det – 0.4 [0,1]

$S \rightarrow NP VP [0.80]$ $NP \rightarrow Det N \quad [0.30]$ $VP \rightarrow V NP \quad [0.20]$

$\begin{array}{c} Det \rightarrow \text{the} \\ Det \rightarrow \text{a} \end{array}$	[0.40] [0.40]			
$V \rightarrow \text{includes}$	[0.05]			
$N \rightarrow \text{meal}$	[0.01]			
$N \rightarrow \text{flight}$	[0.02]		The	flight
		0		2

inc	cludes 3	a meal 4	5	





Det – 0.4 [0,1]

$S \rightarrow NP VP [0.80]$ $VP \rightarrow V NP \quad [0.20]$











$Det \rightarrow \text{the}$	[0.40]			
$Det \rightarrow a$	[0.40]			
$V \rightarrow \text{includes}$	[0.05]			
$N \rightarrow \text{meal}$	[0.01]			
$N \rightarrow \text{flight}$	[0.02]		The	flight
		0		2











$S \rightarrow NP VP$	[0.80]	Det - 0.4	NP		
$NP \rightarrow Det N$	[0.30]				
$VP \rightarrow V NP$	[0.20]				
	L J		N – 0.02		
			[1,2]		
$D_{ot} \rightarrow th_{o}$	[0 /0]	$P = P(NP \to D)$	et N)		
		$P(Det \rightarrow t)$	ne) ·		
$Det \rightarrow a$	[0.40]	$P(N \rightarrow flig)$	ht)		
$V \rightarrow \text{includes}$	[0.05]				
$N \rightarrow \text{meal}$	[0.01]				
$N \rightarrow \text{flight}$	[0.02]	The flight i	ncludes a	meal	
				F	
		0 I 2	3 4	5	









$S \rightarrow NP VP$	[0.80]	Det – 0.4	NP		
$NP \rightarrow Det N$	[0.30]	ΓΟ.11	Г0.21		
$VP \rightarrow V NP$	[0.20]		N – 0.02		
			ГI 21		
			ני, <i>ב</i> ן		
		$P = P(NP \to D)$	et N)		
$Det \rightarrow \text{the}$	[0.40]	$P(Det \rightarrow t)$	ne) ·		
$Det \rightarrow a$	[0.40]	$P(N \rightarrow flig)$	ht)		
$V \rightarrow \text{includes}$	[0.05]				
$N \rightarrow \text{meal}$	[0.01]	P = 0.3 · 0.4 · 0.02 =	= 0.00024		
$N \rightarrow \text{flight}$	[0.02]	The flight i	ncludes a	meal	
		0 1 2	3 4	5	









$S \rightarrow NP VP$	[0.80]		Det – 0.4	NP – 0.0024				
$\frac{NP}{VD} \rightarrow \frac{Det}{V} \frac{N}{N}$	$\begin{bmatrix} 0.30 \end{bmatrix}$		[0,1]	[0,2]				
$VP \rightarrow V NP$	[0.20]			N – 0.02				
				[1,2]				
$Det \rightarrow \text{the}$	[0.40]	P = P	$P(NP \rightarrow D)$	et N)				
$Det \rightarrow a$	[0.40]	F	$(N \rightarrow flig)$	ht)				
$V \rightarrow \text{includes}$ $N \rightarrow \text{meal}$	[0.05] $[0.01]$	$P = 0.3 \cdot 0$	0.4 · 0.02 =	= 0.00024				
$N \rightarrow \text{flight}$	[0.02]	The	flight i	ncludes	a r	neal		
		0	2	3	4	5	5	





Det – 0.4 [0,1]

2

$S \rightarrow NP VP$ $\left[0.80\right]$ $NP \rightarrow Det N \quad [0.30]$ $VP \rightarrow V NP \quad [0.20]$

[0.40] $Det \rightarrow \text{the}$ $Det \rightarrow a \quad [0.40]$ [0.05] $V \rightarrow \text{includes}$ [0.01] $N \rightarrow \text{meal}$ $N \rightarrow \text{flight} [0.02]$ The flight ir

	NP – 0.0024			S – 2.304×10-8	
	[0,2]	[0,3]	[0,4]	[0,5]	
	N – 0.02				
	[1,2]	[1,3]	[1,4]	[1,5]	
		V – 0.05		VP – 1.2×10 ⁻⁵	
		[2,3]	[2,4]	[2,5]	
			Det – 0.4	NP – 0.0012	
			[3,4]	[3,5]	
				N – 0.01	
ncludes a meal [4,5]					
	3	4	5		





Inducing a PCFG







• Use treebank of parsed sentences







- Use treebank of parsed sentences
- To compute probability of a rule, count:

Learning Probabilities





- Use treebank of parsed sentences
- To compute probability of a rule, count:
 - Number of times a nonterminal is expanded:

Learning Probabilities

 $\Sigma_{\gamma} Count(\alpha \rightarrow \gamma)$







- Simplest way:
 - Use treebank of parsed sentences
 - To compute probability of a rule, count:
 - Number of times a nonterminal is expanded:
 - Number of times a nonterminal is expanded by a given rule:

 $\Sigma_{\gamma} Count(\alpha \rightarrow \gamma)$ $Count(\alpha \rightarrow \beta)$







- Use treebank of parsed sentences
- To compute probability of a rule, count:
 - Number of times a nonterminal is expanded:
 - Number of times a nonterminal is expanded by a given rule:

$$P(\alpha \to \beta \,|\, \alpha) = \frac{Count(\alpha)}{\sum_{\gamma} Count(\alpha)}$$

 $\Sigma_{\gamma} Count(\alpha \rightarrow \gamma)$ $Count(\alpha \rightarrow \beta)$

 $\frac{(\alpha \to \beta)}{(\alpha \to \gamma)} = \frac{Count(\alpha \to \beta)}{Count(\alpha)}$







- Use treebank of parsed sentences
- To compute probability of a rule, count:
 - Number of times a nonterminal is expanded:
 - Number of times a nonterminal is expanded by a given rule:

$$P(\alpha \to \beta \,|\, \alpha) = \frac{Count(\alpha)}{\sum_{\gamma} Count(\alpha)}$$

• Alternative: Learn probabilities by re-estimating • (Later)

- $\Sigma_{\gamma} Count(\alpha \rightarrow \gamma)$ $Count(\alpha \rightarrow \beta)$
- $\frac{(\alpha \to \beta)}{(\alpha \to \gamma)} = \frac{Count(\alpha \to \beta)}{Count(\alpha)}$







Probabilistic Parser Development Paradigm

Large

Train

(eg.WSJ 2–21 39,830 sentence Estimate rule probabilities

Usage

Size

	Dev	Test
	Small	Small/Med
, es)	(e.g.WSJ 22)	(e.g.WSJ, 23, 2,416 sentences)
	Tuning/Verification,	Held Out,
	Check for Overfit	Final Evaluation











• Assume a 'gold standard' set of parses for test set







- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?







- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?







- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?
 - Maximally strict: identical to 'gold standard'






- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?
 - Maximally strict: identical to 'gold standard'
 - Partial credit:







- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?
 - Maximally strict: identical to 'gold standard'
 - Partial credit:
 - Constituents in output match those in reference







- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?
 - Maximally strict: identical to 'gold standard'
 - Partial credit:
 - Constituents in output match those in reference
 - Same start point, end point, non-terminal symbol





• How can we compute parse score from constituents?

• Multiple Measures:

Labeled Recall (LR) =

Labeled Precision (LP) =

Parseval

of correct constituents in hypothetical parse

of **total** constituents in **reference** parse

of correct constituents in hypothetical parse

of **total** consituents in **hypothetical** parse





Parseval

• F-measure:

- Combines precision and recall

• Let $\beta \in \mathbb{R}$, $\beta > 0$ that adjusts *P* vs. *R* s.t. $\beta \propto \frac{R}{p}$ • F_{β} -measure is then: $F_{\beta} = (1 + \beta^2) \cdot \frac{P \cdot R}{\beta^2 \cdot P + R}$ • With F1-measure as $F_1 = \frac{2PR}{P+R}$









W UNIVERSITY of WASHINGTON





Evaluation: Example Reference S S NP VP NP S(0,4) NP PP A B A B b С D а b а d С 3 0 2 4



S(0,4)

W UNIVERSITY of WASHINGTON























S(0,4) NP(0, I)VP(1,4)













W UNIVERSITY of WASHINGTON











S(0,4) NP(0, I)VP(1,4)NP(2,4) PP(3,4)

















- Crossing Brackets:
 - siblings:
 - $((A B) C) \{ (0,2), (2,3) \}$ and hyp. has $(A(BC)) - \{ (0,1), (1,3) \}$



TOP

B

• # of constituents where produced parse has bracketings that overlap for the

```
/* crossing is counted based on the brackets */
/* in test rather than gold file (by Mike) */
for(j=0;j<bn2;j++){</pre>
 for(i=0;i<bn1;i++){</pre>
    if(bracket1[i].result != 5 &&
       bracket2[j].result != 5 &&
       ((bracket1[i].start < bracket2[j].start &&</pre>
         bracket1[i].end > bracket2[j].start &&
         bracket1[i].end < bracket2[j].end) ||</pre>
        (bracket1[i].start > bracket2[j].start &&
         bracket1[i].start < bracket2[j].end &&</pre>
         bracket1[i].end > bracket2[j].end))){
```

from evalb.c







State-of-the-Art Parsing

- Parsers trained/tested on Wall Street Journal PTB
 - LR: 94%+;
 - LP: 94%+;
 - Crossing brackets: 1%

- Standard implementation of Parseval:
 - evalb







Evaluation Issues

- Only evaluating constituency
- There are other grammar formalisms:
 - LFG (Constraint-based)
 - Dependency Structure
- Extrinsic evaluation
 - How well does getting the correct parse match the semantics, etc?







Earley Parsing







Earley vs. CKY

- CKY doesn't capture full original structure
 - Can back-convert binarization, terminal conversion
 - Unit non-terminals require change in CKY







Earley vs. CKY

- CKY doesn't capture full original structure
 - Can back-convert binarization, terminal conversion
 - Unit non-terminals require change in CKY
- Earley algorithm
 - Supports parsing efficiently with arbitrary grammars
 - Top-down search
 - Dynamic programming
 - Tabulated partial solutions
 - Some bottom-up constraints







- Another dynamic programming solution
 - Partial parses stored in "chart"
 - Compactly encodes ambiguity
 - $\bullet O(N^3)$
- Chart entries contain:
 - Subtree for a single grammar rule
 - Progress in completing subtree
 - Position of subtree w.r.t. input

Earley Algorithm







- First, left-to-right pass fills out a chart with N+1 states
 - Chart entries sit between words in the input string
 - Keep track of states of the parse at those positions
 - For each word position, chart contains set of states representing all partial parse trees generated so far
 - e.g. chart[0] contains all partial parse trees generated at the beginning of sentence

Earley Algorithm









Chart Entries

- Three types of constituents:
 - Predicted constituents
 - In-progress constituents
 - Completed constituents







- Represented by Dotted Rules
 - Position of indicates type of constituent
- $_0$ Book $_1$ that $_2$ flight $_3$
 - $S \rightarrow \cdot VP$ [0,0] (predicted)
 - $NP \rightarrow Det \cdot Nom$ [1,2] (in progress)
 - $VP \rightarrow VNP$ [0,3] (completed)
- [x,y] tells us what portion of the input is spanned so far by rule
- Each state s_i: <dotted rule>, [<back pointer>, <current position>]

Parse Progress







0 Book 1 that 2 flight 3

- $S \rightarrow VP$, [0,0]
 - First 0 means S constituent begins at the start of input
 - Second 0 means the dot is here too
 - So, this is a top-down prediction







0 Book 1 that 2 flight 3

- $S \rightarrow VP$, [0,0]
 - First 0 means S constituent begins at the start of input
 - Second 0 means the dot is here too
 - So, this is a top-down prediction
- $NP \rightarrow Det \cdot Nom$, [1,2]
 - the NP begins at position 1
 - the dot is at position 2
 - so, Det has been successfully parsed
 - Nom predicted next







0 Book 1 that 2 flight 3 (continued)

- $VP \rightarrow VNP \cdot [0,3]$
 - Successful VP parse of entire input









Successful Parse

- Final answer found by looking at last entry in chart
- If entry resembles $S \rightarrow \alpha \cdot [0,N]$ then input parsed successfully
- Chart will also contain record of all possible parses of input string, given the grammar







Parsing Procedure for the Earley Algorithm

- Move through each set of states in order, applying one of three operations:
 - predictor: add predictions to the chart
 - scanner: read input and add corresponding state to chart
 - **completer**: move dot to right when new constituent found
- Results (new states) added to current or next set of states in chart
- No backtracking and no states removed: keep complete history of parse





function EARLEY-PARSE(words, grammar) returns chart ENQUEUE(($\gamma \rightarrow \bullet S, [0,0]$), chart[θ]) for $i \leftarrow \text{from 0 to LENGTH}(words)$ do for each state in chart[i] do if INCOMPLETE?(*state*) and NEXT-CAT(*state*) is **not** a part of speech **then PREDICTOR**(*state*) elseif INCOMPLETE?(*state*) and NEXT-CAT(*state*) is a part of speech **then SCANNER**(*state*) else **COMPLETER**(*state*) end end return(chart)

Earley Algorithm







procedure **PREDICTOR** $((A \rightarrow a \bullet B \beta, [i,j]))$ for each $(B \rightarrow \gamma)$ in GRAMMAR-RULES-FOR(B, grammar) do ENQUEUE($(B \rightarrow \bullet \gamma, [j,j]), chart[j]$) end

procedure SCANNER($(A \rightarrow a \bullet B \beta, [i, j])$) **if** B ⊂ PARTS-OF-SPEECH(*word*/*j*/) **then** ENQUEUE((B \rightarrow word[j] \bullet , [j,j+1]), chart[j+1])

procedure COMPLETER($(B \rightarrow \gamma \bullet, [j,k])$) for each $(A \rightarrow a \bullet B \beta, [i,j])$ in *chart*[j] do ENQUEUE($(A \rightarrow a B \bullet \beta, [i,k]), chart[k]$) end

Earley Algorithm







3 Main Subroutines of Earley

- Predictor
 - Adds predictions into the chart
- Scanner
- Completer
 - Moves the dot to the right when new constituents are found

• Reads the input words and enters states representing those words into the chart







Predictor

- Intuition:
 - Create new state for top-down prediction of new phrase
- Applied when non part-of-speech non-terminals are to the right of a dot:
 - $S \rightarrow \cdot VP[0,0]$
- Adds new states to current chart
 - One new state for each expansion of the non-terminal in the grammar $VP \rightarrow \cdot V$ [0,0] $VP \rightarrow VNP$ [0,0]









Chart[0]

S0
$$\gamma \rightarrow \cdot S$$
S1 $S \rightarrow \cdot NP VP$ S2 $S \rightarrow \cdot Aux NP VP$ S3 $S \rightarrow \cdot VP$ S4 $NP \rightarrow \cdot Pronoun$ S5 $NP \rightarrow \cdot Proper-Noun$ S6 $NP \rightarrow \cdot Det Nominal$ S7 $VP \rightarrow \cdot Verb$ S8 $VP \rightarrow \cdot Verb NP$ S9 $VP \rightarrow \cdot Verb NP PP$ S10 $VP \rightarrow \cdot Verb PP$ S11 $VP \rightarrow \cdot VP PP$

- [0,0] Dummy start state
- [0.0] Predictor
- [0,0] Predictor
- [0,0] Predictor
- Predictor [0,0] [0,0] Predictor
- [0,0] Predictor
- Predictor [0,0] [0,0] Predictor Predictor [0,0] [0,0] Predictor [0,0] Predictor







Chart[1]

- S12 Verb \rightarrow book \cdot
- S13 $VP \rightarrow Verb \cdot$ $VP \rightarrow Verb \cdot NP$ S14 S15 $VP \rightarrow Verb \cdot NP PP$ S16 $VP \rightarrow Verb \cdot PP$
- S17 $S \rightarrow VP \cdot$

S20

S21

- S18 $VP \rightarrow VP \cdot PP$
- $NP \rightarrow \cdot Pronoun$ S19
 - $NP \rightarrow \cdot Proper-Noun$
 - *NP* → *Det Nominal*
- S22 $PP \rightarrow \cdot Prep NP$

- [0,1] Scanner
- [0,1] Completer
- [0,1] Completer [0,1] Completer
- [0,1] Completer
- [0,1] Completer
- [0,1] Completer
- [1,1] Predictor [1,1] Predictor [1,1] Predictor Predictor [1,1]





S0: $\gamma \rightarrow \cdot S[0,0]$

Book that flight γ • S

W UNIVERSITY of WASHINGTON





SO: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow \cdot VP[0,0]$

Book that flight γ S • VP

W UNIVERSITY of WASHINGTON





SO: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow \cdot VP[0,0]$ S8: $VP \rightarrow \cdot Verb NP[0,0]$








S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow \cdot VP[0,0]$ S8: $VP \rightarrow \cdot Verb NP[0,0]$ S12: Verb $\rightarrow \cdot book$ [0,0]









S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow \cdot VP[0,0]$ S8: $VP \rightarrow \cdot Verb NP[0,0]$ S12: Verb \rightarrow book \cdot [0,1]

> Verb book •









SO: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow \cdot VP[0,0]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$

Verb • book









S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$











S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow \cdot Det Nominal[1,1]$









S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow \cdot Det Nominal[1,1]$ S23: *Det* → • *"that"* [1,1]









S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow \cdot Det Nominal[1,1]$ S23: *Det* → *"that"* • [1,2]









S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow Det \cdot Nominal[1,2]$









S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow Det \cdot Nominal[1,2]$ S25: Nominal \rightarrow • Noun [2,2]







S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow Det \cdot Nominal$ [1,2] S25: *Nominal* → • *Noun* [2,2] S28: Noun \rightarrow "flight" • [2,3]

> Verb book







S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow Det \cdot Nominal[1,2]$ S25: Nominal \rightarrow Noun • [2,3]

> Verb book







S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb \cdot NP[0,1]$ S21: $NP \rightarrow Det Nominal \cdot [1,3]$

> Verb book







S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,1]$ S8: $VP \rightarrow Verb NP \cdot [0,3]$

Verb book







S0: $\gamma \rightarrow \cdot S[0,0]$ S3: $S \rightarrow VP \cdot [0,3]$

Verb book







What About Dead Ends?





S0: $\gamma \rightarrow \cdot S[0,0]$ S1: $S \rightarrow \cdot NP VP [0,0]$

 $NP \rightarrow \bullet Pronoun$ $NP \rightarrow \cdot Proper-Noun$ *NP* → • *Det Nominal*

book

 $\bullet \bullet \bullet$









S0: $\gamma \rightarrow \cdot S[0,0]$ S1: $S \rightarrow \cdot NP VP [0,0]$





 $\bullet \bullet \bullet$



















• We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?







rules like S -> S 'and' S?

• No!

procedure ENQUEUE(*state*, *chart-entry*) if *state* is not already in *chart-entry* then PUSH(*state*, *chart-entry*) end

• We now have a top-down parser in hand. Does it enter infinite loops on







rules like S -> S 'and' S?

• No!

procedure ENQUEUE(*state*, *chart-entry*) if *state* is not already in *chart-entry* then PUSH(state, chart-entry) end

Exercise: parse 'table and chair' using the very simple grammar Nom -> Nom 'and' Nom | 'table' | 'chair'

• We now have a top-down parser in hand. Does it enter infinite loops on





